

## A LARGE-UPDATE INTERIOR POINT ALGORITHM FOR $P_*(\kappa)$ LCP BASED ON A NEW KERNEL FUNCTION

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ABSTRACT. In this paper we generalize large-update primal-dual interior point methods for linear optimization problems in [2] to the  $P_*(\kappa)$  linear complementarity problems based on a new kernel function which includes the kernel function in [2] as a special case. The kernel function is neither self-regular nor eligible. Furthermore, we improve the complexity result in [2] from  $O(\sqrt{n}(\log n)^2 \log \frac{n\mu_0}{\epsilon})$  to  $O(\sqrt{n}(\log n) \log(\log n) \log \frac{n\mu_0}{\epsilon})$ .

### 1. Introduction

In this paper we propose a new large-update interior point algorithm for solving linear complementarity problem(LCP) as follows:

$$s = Mx + q, \quad xs = 0, \quad x \geq 0, \quad s \geq 0, \quad (1)$$

where  $x, s, q \in \mathbf{R}^n$ ,  $M \in \mathbf{R}^{n \times n}$  is a  $P_*(\kappa)$  matrix, and  $xs$  denotes the componentwise product of the vectors  $x$  and  $s$ .

Primal-dual interior point method(IPM) is one of the most efficient numerical methods for various optimization problems. Linear complementarity problems(LCPs) have many applications in science, economics, and engineering([5]).

It is generally agreed that the iteration complexity of the algorithm is an appropriate measure for its efficiency([6]). Most of polynomial-time interior point algorithms are based on the logarithmic barrier function. Peng et al.([11], [12], [13]) proposed a new variant of interior point methods(IPMs) based on self-regular barrier functions and achieved so far the best known complexity result for large-update methods with a specific self-regular barrier function. Roos et al.([1], [2]) proposed new primal-dual IPMs for linear optimization(LO) problems based on eligible barrier functions and proposed the unified scheme for analyzing the algorithm based on four conditions on the kernel function([2]). Cho et al.([3], [4]) extended the algorithm for LO to  $P_*(\kappa)$  LCPs.

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Motivated by their works, we introduce a new class of kernel functions which is the generalized form of the ones in [2] and is not eligible. We obtained  $\mathcal{O}\left(\frac{(1+2\kappa)}{r}n^{\frac{1}{1+r}}(\log n)^{1+r}\log\frac{n\mu_0}{\epsilon}\right)$  iteration complexity for large-update method. Taking  $p = 1$  and  $r = \frac{1+\epsilon}{\log(\log n)}$ , we have  $\mathcal{O}\left((1+2\kappa)\sqrt{n}\log n\log(\log n)\log\frac{n\mu_0}{\epsilon}\right)$  iteration complexity for  $P_*(\kappa)$  LCP which is better than the one in [2].

The paper is organized as follows. In Section 2 we recall the generic IPM and propose some basic concepts for LCP. In Section 3 we introduce a new class of kernel functions and its properties. In Section 4 we derive the complexity result for the algorithm based on a new kernel function.

We will make use of the following notations throughout the paper.  $\mathbf{R}_+^n$  and  $\mathbf{R}_{++}^n$  denote the set of  $n$ -dimensional nonnegative vectors and positive vectors, respectively. For  $x \in \mathbf{R}^n$ ,  $x_{\min}$  denotes the smallest component of the vector  $x$ . We denote  $X$  and  $S$  the diagonal matrices from a vector  $x$  and  $s$ , respectively, i.e.  $X = \text{diag}(x)$  and  $S = \text{diag}(s)$ .  $e$  and  $E$  denote the  $n$ -dimensional vector of ones and the identity matrix, respectively. For  $f(t), g(t) : \mathbf{R}_{++} \rightarrow \mathbf{R}_{++}$ ,  $f(t) = \mathcal{O}(g(t))$  if  $f(t) \leq c_1 g(t)$  for some positive constant  $c_1$  and  $f(t) = \Theta(g(t))$  if  $c_2 g(t) \leq f(t) \leq c_3 g(t)$  for some positive constants  $c_2$  and  $c_3$ .  $I$  denotes the index set, e.g.  $I = \{1, 2, \dots, n\}$ .  $\log$  denotes the natural logarithmic function.

## 2. Preliminaries

In this section, we recall the generic IPM and introduce basic concepts.

**Definition 1.** [8] Let  $\kappa \geq 0$ .  $P_*(\kappa)$  is the class of matrices  $M$  satisfying

$$(1+4\kappa) \sum_{i \in I_+(\xi)} \xi_i [M\xi]_i + \sum_{i \in I_-(\xi)} \xi_i [M\xi]_i \geq 0,$$

where  $\xi \in \mathbf{R}^n$ ,  $[M\xi]_i$  denotes the  $i$ -th component of the vector  $M\xi$  and

$$I_+(\xi) = \{i \in I : \xi_i [M\xi]_i \geq 0\}, \quad I_-(\xi) = \{i \in I : \xi_i [M\xi]_i < 0\}.$$

**Lemma 2.1.** [8] If  $M \in \mathbf{R}^{n \times n}$  is a  $P_*(\kappa)$  matrix, then

$$M' = \begin{pmatrix} -M & E \\ S & X \end{pmatrix}$$

is a nonsingular matrix for any positive diagonal matrices  $X, S \in \mathbf{R}^{n \times n}$ .

**Corollary 2.2.** Let  $M \in \mathbf{R}^{n \times n}$  be a  $P_*(\kappa)$  matrix and  $x, s \in \mathbf{R}_{++}^n$ . Then for all  $c \in \mathbf{R}^n$  the system

$$-M\Delta x + \Delta s = 0, \quad S\Delta x + X\Delta s = c$$

has a unique solution  $(\Delta x, \Delta s)$ .

The basic idea of primal-dual IPMs is to replace the second equation in (1) by the parameterized equation  $xs = \mu e$ ,  $\mu > 0$ . Now we consider the following system:

$$s = Mx + q, \quad Xs = \mu e, \quad x > 0, \quad s > 0. \quad (2)$$

Without loss of generality, we assume that (1) has a strictly feasible point, i.e., there exists  $(x^0, s^0) > 0$  such that  $s^0 = Mx^0 + q$ . For this, the reader refers to [8]. Since  $M$  is a  $P_*(\kappa)$  matrix and (1) is strictly feasible, the system (2) has a unique solution for each  $\mu > 0$ . We denote the solution  $(x(\mu), s(\mu))$  for each  $\mu > 0$ . We call it the  $\mu$ -center. The set of  $\mu$ -centers ( $\mu > 0$ ) is called the central path of (1). The limit of this central path (as  $\mu$  goes to zero) exists and since the limit point satisfies (1), it yields an optimal solution for (1) ([8]). IPMs follow this central path approximately and approach the solution of (1) as  $\mu$  goes to zero.

For given  $(x, s) := (x^0, s^0)$  by applying Newton method to the system (2) we have the following Newton system:

$$-M\Delta x + \Delta s = 0, \quad S\Delta x + X\Delta s = \mu e - xs. \quad (3)$$

By Corollary 2.2, the system (3) has a unique search direction  $(\Delta x, \Delta s)$ . By taking a step along the search direction  $(\Delta x, \Delta s)$ , one constructs a new positive iterate  $(x_+, s_+)$ , where

$$x_+ = x + \alpha\Delta x, \quad s_+ = s + \alpha\Delta s,$$

for some  $\alpha \geq 0$ . To have the motivation of new algorithm we define the following scaled vectors:

$$v := \sqrt{\frac{xs}{\mu}}, \quad d := \sqrt{\frac{x}{s}}, \quad d_x := \frac{v\Delta x}{x}, \quad d_s := \frac{v\Delta s}{s}, \quad (4)$$

whose  $i$ th components are  $\sqrt{x_i s_i / \mu}$ ,  $\sqrt{x_i / s_i}$ ,  $v_i[\Delta x]_i / x_i$ , and  $v_i[\Delta s]_i / s_i$ , respectively. Using (4), we can rewrite the system (3) as follows:

$$-\bar{M}d_x + d_s = 0, \quad d_x + d_s = v^{-1} - v, \quad (5)$$

where  $\bar{M} := DMD$  and  $D := \text{diag}(d)$ . Note that the right side of the second equation in (5) equals the negative gradient of the logarithmic barrier function  $\Psi_l(v)$ , i.e.,

$$d_x + d_s = -\nabla\Psi_l(v), \quad (6)$$

where

$$\Psi_l(v) := \sum_{i=1}^n \psi_l(v_i), \quad \psi_l(t) = \frac{t^2 - 1}{2} - \log t, \quad t > 0.$$

We call  $\psi_l$  the kernel function of the logarithmic barrier function  $\Psi_l(v)$ .

The generic interior point algorithm works as follows. Assume that we are given a strictly feasible point  $(x, s)$  which is in a  $\tau$ -neighborhood of the given  $\mu$ -center. Then we decrease  $\mu$  to  $\mu_+ := (1 - \theta)\mu$ , for some fixed  $\theta \in (0, 1)$  and solve the Newton system (3) to obtain the unique search direction. The positivity condition of a new iterate is ensured with the right choice of the step size  $\alpha$  which is defined by some line search rule. This procedure is repeated until we find a new iterate  $(x_+, s_+)$  that is in a  $\tau$ -neighborhood of the  $\mu_+$ -center and then we let  $\mu := \mu_+$  and  $(x, s) := (x_+, s_+)$ . Then  $\mu$  is again reduced by the

factor  $1 - \theta$  and we solve the Newton system targeting at the new  $\mu_+$ -center, and so on. This process is repeated until  $\mu$  is small enough, say until  $n\mu < \varepsilon$ .

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### Generic Primal-Dual Algorithm

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Input:

a threshold parameter  $\tau > 0$ ;  
 an accuracy parameter  $\varepsilon > 0$ ;  
 a fixed barrier update parameter  $\theta$ ,  $0 < \theta < 1$ ;  
 $(x^0, s^0)$  and  $\mu^0 > 0$  such that  $\Psi_I(x^0, s^0, \mu^0) \leq \tau$ .

begin

$x := x^0$ ;  $s := s^0$ ;  $\mu := \mu^0$ ;

while  $n\mu \geq \varepsilon$  do

begin

$\mu := (1 - \theta)\mu$ ;

while  $\Psi_I(v) > \tau$  do

begin

solve the system (3) for  $\Delta x$  and  $\Delta s$ ;

determine a step size  $\alpha$ ;

$x := x + \alpha\Delta x$ ;

$s := s + \alpha\Delta s$ ;

$v := \sqrt{\frac{xs}{\mu}}$ ;

end

end

end

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When the barrier update parameter  $\theta$  is independent of  $n$ , we call the algorithm a large-update method.

### 3. New kernel function

In this section we define a new class of kernel functions and its properties.

**Definition 2.** The function  $\psi : \mathbf{R}_{++} \rightarrow \mathbf{R}_+$  is called a kernel function if  $\psi$  is twice differentiable and satisfies the following conditions:

$$(a) \psi'(1) = \psi(1) = 0, (b) \psi''(t) > 0, t > 0, (c) \lim_{t \rightarrow 0} \psi(t) = \lim_{t \rightarrow \infty} \psi(t) = \infty.$$

Now we define a new class of kernel functions with parameters  $p$  and  $r$  as follows:

$$\psi(t) := \frac{t^{p+1} - 1}{p + 1} + r(e^{t^{-\frac{1}{r}} - 1} - 1), \quad 0 \leq p \leq 1, \quad 0 < r \leq 1, \quad t > 0. \quad (7)$$

Note that  $\psi(t)$  includes the kernel function defined in [2] as a special case. For  $\psi(t)$  we have the following:

$$\begin{aligned}\psi'(t) &= t^p - t^{-\frac{1}{r}-1}e^{t^{-\frac{1}{r}-1}}, \\ \psi''(t) &= pt^{p-1} + \left(\frac{1}{r} + \left(\frac{1}{r} + 1\right)t^{\frac{1}{r}}\right)t^{-\frac{2}{r}-2}e^{t^{-\frac{1}{r}-1}}, \\ \psi'''(t) &= p(p-1)t^{p-2} - \left(\frac{1}{r^2} + \frac{3}{r}\left(\frac{1}{r} + 1\right)t^{\frac{1}{r}} + \left(\frac{1}{r} + 1\right)\left(\frac{1}{r} + 2\right)t^{\frac{2}{r}}\right)t^{-\frac{3}{r}-3}e^{t^{-\frac{1}{r}-1}.\end{aligned}\tag{8}$$

In this paper, we replace the function  $\Psi_l(v)$  in (6) with the function  $\Psi(v)$  as follows:

$$d_x + d_s = -\nabla\Psi(v),\tag{9}$$

where  $\Psi(v) = \sum_{i=1}^n \psi(v_i)$  and  $\psi(t)$  is defined in (7). Hence the new search direction  $(\Delta x, \Delta s)$  is obtained by solving the following modified Newton-system:

$$-M\Delta x + \Delta s = 0, \quad S\Delta x + X\Delta s = -\mu v \nabla\Psi(v).\tag{10}$$

Since  $\Psi(v)$  is strictly convex and minimal at  $v = e$ , we have

$$\Psi(v) = 0 \Leftrightarrow v = e \Leftrightarrow x = x(\mu), \quad s = s(\mu).$$

We use  $\Psi(v)$  as the proximity function to measure the distance between the current iterate and the  $\mu$ -center. Also, we define the norm-based proximity measure  $\delta(v)$  as follows:

$$\delta(v) := \frac{1}{2} \|\nabla\Psi(v)\| = \frac{1}{2} \|d_x + d_s\|.\tag{11}$$

In the following we give properties of  $\psi(t)$  which are essential to the complexity analysis.

**Lemma 3.1.** *Let  $\psi(t)$  be as defined in (7). Then we have the following:*

- (i)  $\psi(t)$  is exponentially convex,  $t > 0$ .
- (ii)  $\psi''(t)$  is monotonically decreasing,  $t > 0$ .

*Proof.* For (i), by Lemma 2.1.2 in [13], it suffices to show the function  $\psi(t)$  satisfies  $t\psi''(t) + \psi'(t) \geq 0$  for all  $t > 0$ . Using (8), we have

$$t\psi''(t) + \psi'(t) = (p+1)t^p + \frac{1}{r}t^{-\frac{2}{r}-1}e^{t^{-\frac{1}{r}-1}} + \frac{1}{r}t^{-\frac{1}{r}-1}e^{t^{-\frac{1}{r}-1}} \geq 0, \quad t > 0.$$

For (ii), from (8),  $\psi'''(t) < 0$ . This completes the proof.  $\square$

*Remark 1.* Recall that the function  $\psi : \mathbf{R}_{++} \rightarrow \mathbf{R}_+$  is eligible if  $\psi$  is three times differentiable and satisfies the following conditions([2]):

- (a)  $t\psi''(t) + \psi'(t) > 0$ ,  $t < 1$ ,
- (b)  $t\psi''(t) - \psi'(t) > 0$ ,  $t > 1$ ,
- (c)  $\psi'''(t) < 0$ ,  $t > 0$ ,
- (d)  $2\psi''(t)^2 - \psi'(t)\psi'''(t) > 0$ ,  $t < 1$ .

Using (8), we have

$$t\psi''(t) - \psi'(t) = (p-1)t^p + \frac{1}{r}t^{-\frac{2}{r}-1}e^{t^{-\frac{1}{r}-1}} + \left(\frac{1}{r} + 2\right)t^{-\frac{1}{r}-1}e^{t^{-\frac{1}{r}-1}}.$$

Since

$$t\psi''(t) - \psi'(t) < -7 < 0 \text{ for } p = \frac{1}{2}, r = \frac{1}{4}, t = 2^8,$$

condition (b) is not satisfied. Hence  $\psi(t)$  is not eligible. Note that the kernel function in [2] is eligible.

**Lemma 3.2.** *For  $\psi(t)$  and  $p \in [0, 1]$ , we have*

$$\frac{1}{p+1} \sum_{i=1}^n v_i^{p+1} \leq \Psi(v) + \frac{(pr+r+1)n}{p+1}.$$

*Proof.* Since  $re^{t^{-\frac{1}{r}-1}} > 0$ , we have

$$\psi(t) = \frac{t^{p+1}}{p+1} - \frac{1}{p+1} + re^{t^{-\frac{1}{r}-1}} - r \geq \frac{t^{p+1}}{p+1} - \frac{1}{p+1} - r.$$

So we have  $\frac{t^{p+1}}{p+1} \leq \psi(t) + \frac{pr+r+1}{p+1}$ . Hence we have

$$\frac{1}{p+1} \sum_{i=1}^n v_i^{p+1} \leq \Psi(v) + \frac{(pr+r+1)n}{p+1}.$$

This completes the proof.  $\square$

Define  $\psi_b(t) := r(e^{t^{-\frac{1}{r}-1}} - 1)$ . Then we have  $\psi(t) := \frac{t^{p+1}-1}{p+1} + \psi_b(t)$ . Since  $\psi'_b(t) = -t^{-\frac{1}{r}-1}e^{t^{-\frac{1}{r}-1}} < 0$ ,  $\psi_b(t)$  is monotonically decreasing in  $t$ .

**Lemma 3.3.** *Let  $\beta \geq 1$ . Then  $\psi(\beta t) \leq \psi(t) + \frac{t^{p+1}}{p+1}(\beta^{p+1} - 1)$ .*

*Proof.* Since  $\psi_b(t)$  is monotonically decreasing in  $t$ ,  $\psi_b(\beta t) - \psi_b(t) \leq 0$  for  $\beta \geq 1$ . Hence we have

$$\begin{aligned} \psi(\beta t) &= \frac{(\beta t)^{p+1} - 1}{p+1} + \psi_b(\beta t) \\ &= \frac{t^{p+1} - 1}{p+1} + \psi_b(t) + \frac{1}{p+1}(\beta^{p+1}t^{p+1} - t^{p+1}) + \psi_b(\beta t) - \psi_b(t) \\ &= \psi(t) + \frac{t^{p+1}}{p+1}(\beta^{p+1} - 1) + \psi_b(\beta t) - \psi_b(t) \\ &\leq \psi(t) + \frac{t^{p+1}}{p+1}(\beta^{p+1} - 1). \end{aligned}$$

This completes the proof.  $\square$

In the following we obtain an estimate for the effect of a  $\mu$ -update on the value of  $\Psi(v)$ .

**Theorem 3.4.** *Let  $0 \leq \theta < 1$  and  $v_+ = \frac{v}{\sqrt{1-\theta}}$ . Then we have*

$$\Psi(v_+) \leq \Psi(v) + \frac{\theta}{(1-\theta)^{\frac{1+p}{2}}} \left( \Psi(v) + \frac{(pr+r+1)n}{p+1} \right).$$

*Proof.* Using Lemma 3.3 with  $\beta = \frac{1}{\sqrt{1-\theta}}$  and Lemma 3.2, we have

$$\begin{aligned} \Psi(v_+) &= \Psi(\beta v) = \sum_{i=1}^n \psi(\beta v_i) \leq \sum_{i=1}^n \left( \psi(v_i) + \frac{1}{p+1} (\beta^{p+1} - 1) v_i^{p+1} \right) \\ &= \Psi(v) + \left( \frac{1}{(1-\theta)^{\frac{1+p}{2}}} - 1 \right) \frac{1}{p+1} \sum_{i=1}^n v_i^{p+1} \\ &\leq \Psi(v) + \frac{1 - (1-\theta)^{\frac{1+p}{2}}}{(1-\theta)^{\frac{1+p}{2}}} \left( \Psi(v) + \frac{(pr+r+1)n}{p+1} \right). \end{aligned}$$

Since  $1 - (1-\theta)^{\frac{1+p}{2}} \leq \theta$  for  $0 \leq \theta < 1$ ,

$$\Psi(v_+) \leq \Psi(v) + \frac{\theta}{(1-\theta)^{\frac{1+p}{2}}} \left( \Psi(v) + \frac{(pr+r+1)n}{p+1} \right).$$

This completes the proof.  $\square$

Note that at the start of outer iteration of the algorithm, i.e., just before the update of  $\mu$  with the factor  $1-\theta$ , we have  $\Psi(v) \leq \tau$ . During the inner iteration we have

$$\begin{aligned} \Psi(v_+) &\leq \Psi(v) + \frac{\theta}{(1-\theta)^{\frac{1+p}{2}}} \left( \Psi(v) + \frac{(pr+r+1)n}{p+1} \right) \\ &\leq \tau + \frac{\theta}{(1-\theta)^{\frac{1+p}{2}}} \left( \tau + \frac{(pr+r+1)n}{p+1} \right). \end{aligned}$$

Each subsequent inner iteration will rise to a decrease of the value of  $\Psi(v)$ . Denote

$$\tilde{\Psi}_0 := \tau + \frac{\theta}{(1-\theta)^{\frac{1+p}{2}}} \left( \tau + \frac{(pr+r+1)n}{p+1} \right). \quad (12)$$

Define  $\Psi_0$  the value of  $\Psi(v)$  after the  $\mu$ -update. Then  $\Psi_0 \leq \tilde{\Psi}_0$ .

**Lemma 3.5.** *Define  $\varrho : [0, \infty) \rightarrow [1, \infty)$  be the inverse function of  $\psi(t)$  for  $t \geq 1$ . For  $0 \leq p \leq 1$  and  $u \geq 0$  we have*

$$\varrho(u) \geq (1 + (p+1)u)^{\frac{1}{1+p}}.$$

*Proof.* Let  $u = \psi(t)$ ,  $t \geq 1$ . Since  $\psi_b(t)$  is monotonically decreasing in  $t$  and  $\psi_b(1) = 0$ ,  $\psi_b(t) < 0$  for  $t > 1$ . Hence  $u = \psi(t) = \frac{t^{p+1}-1}{p+1} + \psi_b(t) \leq \frac{t^{p+1}-1}{p+1}$ ,  $t \geq 1$ . This implies  $(p+1)u + 1 \leq t^{p+1}$ . By the definition of  $\varrho$ ,  $\varrho(u) = t \geq (1 + (p+1)u)^{\frac{1}{1+p}}$ . This completes the proof.  $\square$

From Lemma 3.1 (ii), we cite the following lemma in [2] without proof.

**Lemma 3.6.** (Theorem 4.9 in [2]) *Let  $\delta(v)$  be as defined in (11). Then we have*

$$\delta(v) \geq \frac{1}{2} \psi'(\varrho(\Psi(v))).$$

For notational convenience we denote  $\delta := \delta(v)$  and  $\Psi := \Psi(v)$ .

**Lemma 3.7.** *Let  $\delta$  be as defined in (11). Then for all  $\Psi \geq 1$  and  $0 \leq p \leq 1$  we have*

$$\delta \geq \frac{1}{4} ((p+1)\Psi)^{\frac{p}{1+p}}.$$

*Proof.* By Lemma 3.6, Lemma 3.5, and  $\psi''(t) > 0$ ,

$$\begin{aligned} \delta &\geq \frac{1}{2} \psi'(\varrho(\Psi)) \geq \frac{1}{2} \psi'((1+(p+1)\Psi)^{\frac{1}{1+p}}) \\ &= \frac{1}{2} \left( (1+(p+1)\Psi)^{\frac{p}{1+p}} - e^{(1+(p+1)\Psi)^{-\frac{1}{r(1+p)}-1}} \frac{1}{(1+(p+1)\Psi)^{\frac{1+r}{r(1+p)}}} \right) \\ &\geq \frac{1}{2} \left( (1+(p+1)\Psi)^{\frac{p}{1+p}} - \frac{1}{(1+(p+1)\Psi)^{\frac{1+r}{r(1+p)}}} \right) \\ &\geq \frac{1}{2} \left( (1+(p+1)\Psi)^{\frac{p}{1+p}} - \frac{1}{(1+(p+1)\Psi)^{\frac{1}{1+p}}} \right) \\ &= \frac{1}{2} \frac{(p+1)\Psi}{(1+(p+1)\Psi)^{\frac{1}{1+p}}} \geq \frac{1}{4} ((p+1)\Psi)^{\frac{p}{1+p}}, \end{aligned}$$

where the third inequality is satisfied from  $e^{(1+(p+1)\Psi)^{-\frac{1}{r(1+p)}-1}} \leq 1$  and the last inequality from the fact  $1 \leq (p+1)\Psi$ . This completes the proof.  $\square$

#### 4. Complexity result

In this section we compute a feasible step size and the decrease of the proximity function during an inner iteration and give the complexity result of the algorithm. For fixed  $\mu$  if we take a step size  $\alpha$ , then we have new iterates  $x_+ = x + \alpha \Delta x$ ,  $s_+ = s + \alpha \Delta s$ . Using (4), we have

$$x_+ = x \left( e + \alpha \frac{\Delta x}{x} \right) = x \left( e + \alpha \frac{d_x}{v} \right) = \frac{x}{v} (v + \alpha d_x)$$

and

$$s_+ = s \left( e + \alpha \frac{\Delta s}{s} \right) = s \left( e + \alpha \frac{d_s}{v} \right) = \frac{s}{v} (v + \alpha d_s).$$

Thus we have

$$v_+ = \sqrt{\frac{x_+ s_+}{\mu}} = \sqrt{(v + \alpha d_x)(v + \alpha d_s)}.$$



Define for  $\alpha > 0$ ,  $f(\alpha) = \Psi(v_+) - \Psi(v)$ . Then  $f(\alpha)$  is the difference between proximities of a new iterate and a current iterate for fixed  $\mu$ . Using Lemma 3.1 (i), we have

$$\Psi(v_+) = \Psi(\sqrt{(v + \alpha d_x)(v + \alpha d_s)}) \leq \frac{1}{2}(\Psi(v + \alpha d_x) + \Psi(v + \alpha d_s)).$$

Hence we have  $f(\alpha) \leq f_1(\alpha)$ , where

$$f_1(\alpha) := \frac{1}{2}(\Psi(v + \alpha d_x) + \Psi(v + \alpha d_s)) - \Psi(v).$$

We have  $f(0) = f_1(0) = 0$ . Taking the derivative of  $f_1(\alpha)$  with respect to  $\alpha$ , we have

$$f_1'(\alpha) = \frac{1}{2} \sum_{i=1}^n (\psi'(v_i + \alpha [d_x]_i) [d_x]_i + \psi'(v_i + \alpha [d_s]_i) [d_s]_i),$$

where  $[d_x]_i$  and  $[d_s]_i$  denote the  $i$ -th components of the vectors  $d_x$  and  $d_s$ , respectively. Using (9) and (11), we have

$$f_1'(0) = \frac{1}{2} \nabla \Psi(v)^T (d_x + d_s) = -\frac{1}{2} \nabla \Psi(v)^T \nabla \Psi(v) = -2\delta(v)^2.$$

Differentiating  $f_1'(\alpha)$  with respect to  $\alpha$ , we have

$$f_1''(\alpha) = \frac{1}{2} \sum_{i=1}^n (\psi''(v_i + \alpha [d_x]_i) [d_x]_i^2 + \psi''(v_i + \alpha [d_s]_i) [d_s]_i^2).$$

Since  $f_1''(\alpha) > 0$ ,  $f_1(\alpha)$  is strictly convex in  $\alpha$  unless  $d_x = d_s = 0$ . Since  $M$  is a  $P_*(\kappa)$  matrix and  $M\Delta x = \Delta s$  from (10), for  $\Delta x \in \mathbf{R}^n$ ,

$$(1 + 4\kappa) \sum_{i \in I_+} [\Delta x]_i [\Delta s]_i + \sum_{i \in I_-} [\Delta x]_i [\Delta s]_i \geq 0,$$

where  $I_+ = \{i \in I : [\Delta x]_i [\Delta s]_i \geq 0\}$ ,  $I_- = I - I_+$ . Since  $d_x d_s = \frac{v^2 \Delta x \Delta s}{x s} = \frac{\Delta x \Delta s}{\mu}$  and  $\mu > 0$ , we have

$$(1 + 4\kappa) \sum_{i \in I_+} [d_x]_i [d_s]_i + \sum_{i \in I_-} [d_x]_i [d_s]_i \geq 0.$$

For convenience we denote  $\sigma_+ := \sum_{i \in I_+} [d_x]_i [d_s]_i$  and  $\sigma_- := -\sum_{i \in I_-} [d_x]_i [d_s]_i$ . In the following we cite some lemmas in [4] without proof.

**Lemma 4.1.** (Modification of Lemma 4.1 in [4])  $\sigma_+ \leq \delta^2$  and  $\sigma_- \leq (1 + 4\kappa)\delta^2$ .

**Lemma 4.2.** (Modification of Lemma 4.2 in [4])  $\sum_{i=1}^n ([d_x]_i^2 + [d_s]_i^2) \leq 4(1 + 2\kappa)\delta^2$ ,  $\|d_x\| \leq 2\delta\sqrt{1 + 2\kappa}$ , and  $\|d_s\| \leq 2\delta\sqrt{1 + 2\kappa}$ .

**Lemma 4.3.** (Modification of lemma 4.3 in [4])  $f_1''(\alpha) \leq 2(1 + 2\kappa) \delta^2 \psi''(v_{min} - 2\alpha\delta\sqrt{1 + 2\kappa})$ .

**Lemma 4.4.** (Modification of lemma 4.4 in [4])  $f_1'(\alpha) \leq 0$  if  $\alpha$  is satisfying

$$(13) \quad -\psi'(v_{min} - 2\alpha\delta\sqrt{1+2\kappa}) + \psi'(v_{min}) \leq \frac{2\delta}{\sqrt{1+2\kappa}}.$$

**Lemma 4.5.** (Modification of lemma 4.5 in [4]) Define  $\rho : [0, \infty) \rightarrow (0, 1]$  be the inverse function of  $-\frac{1}{2}\psi'(t)$  for  $0 < t \leq 1$  and  $a := 1 + \frac{1}{\sqrt{1+2\kappa}}$ . Then the largest step size  $\alpha$  satisfying (13) is given by

$$\hat{\alpha} := \frac{1}{2\delta\sqrt{1+2\kappa}} (\rho(\delta) - \rho(a\delta)).$$

**Lemma 4.6.** (Modification of lemma 4.6 in [4]) Let  $\rho$  and  $\hat{\alpha}$  be as defined in Lemma 4.5. Then

$$\hat{\alpha} \geq \frac{1}{(1+2\kappa)\psi''(\rho(a\delta))}.$$

Define

$$\bar{\alpha} := \frac{1}{(1+2\kappa)\psi''(\rho(a\delta))}. \quad (14)$$

Then we have  $\bar{\alpha} \leq \hat{\alpha}$ .

**Lemma 4.7.** Let  $\bar{\alpha}$  be as defined in (14). Then for  $a = 1 + \frac{1}{\sqrt{1+2\kappa}}$  and  $\kappa \geq 0$ , we have

$$\bar{\alpha} \geq \frac{1}{(1+2\kappa) \left( p + (2a\delta + 1) \left( \frac{2}{r} + 1 \right) (1 + \log(2a\delta + 1))^{1+r} \right)}.$$

*Proof.* Using the definition of  $\rho$ , we have  $-\frac{1}{2}\psi'(\rho(a\delta)) = a\delta$ . Let  $z = \rho(a\delta)$ . Then  $-\psi'(z) = 2a\delta$  and  $0 < z \leq 1$ . From (8), we have  $-z^p + z^{-\frac{1}{r}-1}e^{z^{-\frac{1}{r}-1}} = 2a\delta$ . Then for  $0 < z \leq 1$ ,

$$z^{-\frac{1}{r}-1}e^{z^{-\frac{1}{r}-1}} = 2a\delta + z^p \leq 2a\delta + 1. \quad (15)$$

By taking the natural logarithmic function on both sides of (15), we have

$$z^{-\frac{1}{r}} - 1 - \left( \frac{1}{r} + 1 \right) \log z \leq \log(2a\delta + 1). \quad (16)$$

Using (16) and  $0 < z \leq 1$ , we obtain

$$z^{-\frac{1}{r}} \leq 1 + \log(2a\delta + 1) + \left( \frac{1}{r} + 1 \right) \log z \leq 1 + \log(2a\delta + 1).$$

This implies

$$\begin{aligned} z^{p-1} &\leq (1 + \log(2a\delta + 1))^{r(1-p)} \leq (1 + \log(2a\delta + 1))^{1+r}, \\ z^{-1} &\leq (1 + \log(2a\delta + 1))^r \leq (1 + \log(2a\delta + 1))^{1+r}, \\ z^{-\frac{1}{r}-1} &\leq (1 + \log(2a\delta + 1))^{1+r}. \end{aligned} \quad (17)$$

From (14), we have for  $0 < z \leq 1$  and  $0 \leq p \leq 1$ ,

$$\begin{aligned}
\bar{\alpha} &= \frac{1}{(1+2\kappa)\psi''(\rho(a\delta))} = \frac{1}{(1+2\kappa)\psi''(z)} \\
&= \frac{1}{(1+2\kappa)\left(pz^{p-1} + \frac{1}{r}z^{-\frac{2}{r}-2}e^{z^{-\frac{1}{r}-1}} + \left(\frac{1}{r}+1\right)z^{-\frac{1}{r}-2}e^{z^{-\frac{1}{r}-1}}\right)} \\
&\geq \frac{1}{(1+2\kappa)\left(pz^{p-1} + \frac{1}{r}(2a\delta+1)z^{-\frac{1}{r}-1} + \left(\frac{1}{r}+1\right)(2a\delta+1)z^{-1}\right)} \\
&= \frac{1}{(1+2\kappa)\left(pz^{p-1} + (2a\delta+1)\left(\frac{1}{r}z^{-\frac{1}{r}-1} + \left(\frac{1}{r}+1\right)z^{-1}\right)\right)} \\
&\geq \frac{1}{(1+2\kappa)\left(p + (2a\delta+1)\left(\frac{2}{r}+1\right)\right)(1+\log(2a\delta+1))^{1+r}},
\end{aligned}$$

where the first inequality follows from (15) and the last inequality from (17). This proves the lemma.  $\square$

Define

$$\tilde{\alpha} = \frac{1}{(1+2\kappa)\left(p + (2a\delta+1)\left(\frac{2}{r}+1\right)\right)(1+\log(2a\delta+1))^{1+r}}. \quad (18)$$

Note that  $\tilde{\alpha} \leq \bar{\alpha}$ . We will use  $\tilde{\alpha}$  as the default step size.

**Lemma 4.8.** (Lemma 1.3.3 in [13]) *Suppose that  $h(t)$  is a twice differentiable convex function with  $h(0) = 0$  and  $h'(0) < 0$  and  $h(t)$  attains its global minimum at  $t^* > 0$  and  $h''(t)$  is increasing with respect to  $t$ . Then for any  $t \in [0, t^*]$ ,*

$$h(t) \leq \frac{th'(0)}{2}.$$

**Lemma 4.9.** (Modification of lemma 4.8 in [4]) *If the step size  $\alpha$  is such that  $\alpha \leq \tilde{\alpha}$ , then*

$$f(\alpha) \leq -\alpha\delta^2.$$

In our algorithm we assume that  $\tau \geq 1$ . Using Lemma 3.7 and the fact  $\Psi \geq \tau$ , we have

$$\delta \geq \frac{1}{4}\left((p+1)\Psi\right)^{\frac{p}{1+p}} \geq \frac{1}{4}. \quad (19)$$

**Lemma 4.10.** *For  $0 < r \leq 1$  the function*

$$g(\delta) = -\frac{\delta}{(1+\log(2a\delta+1))^{1+r}}$$

*is monotonically decreasing in  $\delta$ .*

*Proof.* It suffices to show that the function  $-g(\delta)$  is monotonically increasing in  $\delta$ . If we differentiate  $-g(\delta)$  with respect to  $\delta$ , we have

$$-g'(\delta) = \frac{(2a\delta + 1)(1 + \log(2a\delta + 1))^{1+r} - 2a\delta(1+r)(1 + \log(2a\delta + 1))^r}{(2a\delta + 1)(1 + \log(2a\delta + 1))^{2(1+r)}}.$$

Since the denominator is strictly positive, it is enough to show that numerator is positive. The numerator is

$$\begin{aligned} & (2a\delta + 1)(1 + \log(2a\delta + 1))^{1+r} - 2a\delta(1+r)(1 + \log(2a\delta + 1))^r \\ &= (1 + \log(2a\delta + 1))^r ((1 + \log(2a\delta + 1))(2a\delta + 1) - 2a\delta(1+r)) \\ &= (1 + \log(2a\delta + 1))^r (1 + (2a\delta + 1)\log(2a\delta + 1) - 2a\delta r). \end{aligned}$$

Let  $\tilde{g}(\delta) := 1 + (2a\delta + 1)\log(2a\delta + 1) - 2a\delta r$ . Then  $\tilde{g}'(\delta) = 2a\log(2a\delta + 1) + 2a - 2ar$ . From  $1 < a \leq 2$  and (19), we have  $\tilde{g}'(\delta) > 0$ . Since  $\tilde{g}(\frac{1}{4}) > 0$ ,  $\tilde{g}(\delta) > 0$  for  $\delta \geq \frac{1}{4}$ . Hence,  $-g(\delta)$  is monotonically increasing in  $\delta$ . This completes the proof.  $\square$

**Theorem 4.11.** *Let  $\tilde{\alpha}$  be as defined in (18). Then*

$$f(\tilde{\alpha}) \leq -\frac{((p+1)\Psi)^{\frac{p}{1+p}}}{16(1+2\kappa)(p+\frac{4}{r}+2)\left(1+\log\left(\frac{a}{2}((p+1)\Psi_0)^{\frac{p}{1+p}}+1\right)\right)^{1+r}}.$$

*Proof.* Using (18) and Lemma 4.9, we have

$$\begin{aligned} f(\tilde{\alpha}) &\leq -\frac{\delta^2}{(1+2\kappa)(p+(2a\delta+1)(\frac{2}{r}+1))(1+\log(2a\delta+1))^{1+r}} \\ &\leq -\frac{\delta^2}{(1+2\kappa)(4p\delta+\delta(2a+4)(\frac{2}{r}+1))(1+\log(2a\delta+1))^{1+r}} \\ &= -\frac{\delta}{2(1+2\kappa)(2p+\frac{2a}{r}+\frac{4}{r}+a+2)(1+\log(2a\delta+1))^{1+r}} \\ &\leq -\frac{\delta}{4(1+2\kappa)(p+\frac{4}{r}+2)(1+\log(2a\delta+1))^{1+r}} \\ &\leq -\frac{\frac{1}{4}((p+1)\Psi)^{\frac{p}{1+p}}}{4(1+2\kappa)(p+\frac{4}{r}+2)\left(1+\log\left(\frac{a}{2}((p+1)\Psi)^{\frac{p}{1+p}}+1\right)\right)^{1+r}} \\ &= -\frac{((p+1)\Psi)^{\frac{p}{1+p}}}{16(1+2\kappa)(p+\frac{4}{r}+2)\left(1+\log\left(\frac{a}{2}((p+1)\Psi)^{\frac{p}{1+p}}+1\right)\right)^{1+r}} \\ &\leq -\frac{((p+1)\Psi)^{\frac{p}{1+p}}}{16(1+2\kappa)(p+\frac{4}{r}+2)\left(1+\log\left(\frac{a}{2}((p+1)\Psi_0)^{\frac{p}{1+p}}+1\right)\right)^{1+r}}, \end{aligned}$$

where the second inequality is satisfied from (19), third inequality from  $1 < a \leq 2$ , the fourth inequality from Lemma 3.7 and Lemma 4.10, and the last inequality from the definition of  $\Psi_0$ . This completes the proof.  $\square$

**Lemma 4.12.** (Lemma 1.3.2 in [13]) *Let  $t_0, t_1, \dots, t_J$  be a sequence of positive numbers such that*

$$t_{j+1} \leq t_j - \gamma t_j^{1-\lambda}, \quad j = 0, 1, \dots, J-1,$$

where  $\gamma > 0$  and  $0 < \lambda \leq 1$ . Then  $J \leq \lfloor \frac{t_0^\lambda}{\gamma\lambda} \rfloor$ .

We define the value of  $\Psi(v)$  after the  $\mu$ -update as  $\Psi_0$  and the subsequent values in the same outer iteration  $\Psi_k$ ,  $k = 1, 2, \dots$ . Let  $K$  denote the total number of inner iterations in the outer iteration. Then we have

$$\Psi_{K-1} > \tau, \quad 0 \leq \Psi_K \leq \tau.$$

**Lemma 4.13.** *Let  $\tilde{\Psi}_0$  be as defined in (12) and  $K$  be the total number of inner iterations in the outer iteration. Then we have*

$$K \leq 16(1+2\kappa) \left( p + \frac{4}{r} + 2 \right) (p+1)^{\frac{1}{1+p}} \left( 1 + \log((p+1)\tilde{\Psi}_0) \right)^{1+r} \tilde{\Psi}_0^{\frac{1}{1+p}}.$$

*Proof.* By Theorem 4.11 with  $\gamma = \frac{(p+1)^{\frac{p}{1+p}}}{16(1+2\kappa) \left( p + \frac{4}{r} + 2 \right) \left( 1 + \log\left(\frac{a}{2}((p+1)\Psi_0)^{\frac{p}{1+p}} + 1\right) \right)^{1+r}}$

and  $\lambda = \frac{1}{1+p}$ , we have

$$K \leq \frac{16(1+2\kappa) \left( p + \frac{4}{r} + 2 \right) \left( 1 + \log\left(\frac{a}{2}((p+1)\Psi_0)^{\frac{p}{1+p}} + 1\right) \right)^{1+r}}{(p+1)^{\frac{p}{1+p}}} (p+1)\Psi_0^{\frac{1}{1+p}}.$$

Since  $\Psi_0 \leq \tilde{\Psi}_0$  and  $1 < a \leq 2$ , we have the result.  $\square$

**Theorem 4.14.** *Let a  $P_*(\kappa)$  LCP be given and  $\tau \geq 1$ . Then the total number of iterations to have an approximate solution with  $n\mu < \epsilon$  is bounded by*

$$\left\lceil 16(1+2\kappa) \left( p + \frac{4}{r} + 2 \right) (p+1)^{\frac{1}{1+p}} \left( 1 + \log((p+1)\tilde{\Psi}_0) \right)^{1+r} \tilde{\Psi}_0^{\frac{1}{1+p}} \right\rceil \cdot \left\lceil \frac{1}{\theta} \log \frac{n\mu_0}{\epsilon} \right\rceil,$$

where  $\epsilon > 0$  is the desired accuracy,  $\mu_0 > 0$  is given, and  $\theta$ ,  $0 < \theta < 1$ , is the given barrier update parameter.

*Proof.* If the central path parameter  $\mu$  has the initial value  $\mu_0 > 0$  and is updated by multiplying  $1 - \theta$  with  $0 < \theta < 1$ , then after at most

$$\left\lceil \frac{1}{\theta} \log \frac{n\mu_0}{\epsilon} \right\rceil$$

iterations we have  $n\mu < \epsilon([14])$ . For the total number of iterations, we multiply the number of inner iterations by that of outer iterations. i.e.,

$$\left[ 16(1+2\kappa) \left( p + \frac{4}{r} + 2 \right) (p+1)^{\frac{1}{1+p}} (1 + \log((p+1)\tilde{\Psi}_0))^{1+r} \tilde{\Psi}_0^{\frac{1}{1+p}} \right] \cdot \left\lceil \frac{1}{\theta} \log \frac{n\mu_0}{\epsilon} \right\rceil.$$

This completes the proof.  $\square$

*Remark 2.* Taking  $\tau = \mathcal{O}(n)$  and  $\theta = \Theta(1)$ , the large-update algorithm has

$$\mathcal{O} \left( \frac{(1+2\kappa)}{r} n^{\frac{1}{1+p}} (\log n)^{1+r} \log \frac{n\mu_0}{\epsilon} \right)$$

iteration complexity. In particular, for  $r = \frac{1+\epsilon}{\log(\log n)}$  with a sufficiently small  $\epsilon > 0$ , we have  $\frac{1}{r} (\log n)^{1+r} = \frac{\epsilon^{1+\epsilon}}{1+\epsilon} (\log n) \log(\log n)$ . So we have  $\mathcal{O}((1+2\kappa)\sqrt{n}(\log n) \log(\log n) \log \frac{n\mu_0}{\epsilon})$  iteration complexity with  $p = 1$  and  $r = \frac{1+\epsilon}{\log(\log n)}$ . This complexity result improves the one in [2].

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