

## SIMULATIVE AND MUTANT WFI-ALGEBRAS

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**Abstract.** The notion of simulative and mutant WFI-algebras is introduced, and several properties are investigated. Characterizations of a simulative WFI-algebra are established. A relation between an associative WFI-algebra and a simulative WFI-algebra is given. Some types for a simulative WFI-algebra to be mutant are found.

### 1. Introduction

In 1990, W. M. Wu [3] introduced the notion of fuzzy implication algebras (FI-algebra, for short), and investigated several properties. In [2], Z. Li and C. Zheng introduced the notion of distributive (resp. regular, commutative) FI-algebras, and investigated the relations between such FI-algebras and MV-algebras. In [1], Y. B. Jun discussed WFI-algebras (weak fuzzy implication algebras) which are weaker than FI-algebras, and gave a characterization of a WFI-algebra. He introduced the notion of associative (resp. normal, medial) WFI-algebras, and investigated several properties. He gave conditions for a WFI-algebra to be associative/medial, and provided characterizations of associative/medial WFI-algebras, and showed that every associative WFI-algebra is a group in which every element is an involution. He also verified that the class

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of all medial WFI-algebras is a variety. In this paper, we introduce the notion of simulative and mutant WFI-algebras and investigate some properties. We establish characterizations of a simulative WFI-algebra. We give a relation between an associative WFI-algebra and a simulative WFI-algebra. We find some types for a simulative WFI-algebra to be mutant.

## 2. Preliminaries

We investigate an algebra  $\mathcal{G} := (G; \ominus, 1)$ , consisting of a set  $G$ , a binary operation  $\ominus$  and a special element  $1$ , which satisfies the four axioms:

- (a1)  $x \ominus (y \ominus z) = y \ominus (x \ominus z)$ ,
- (a2)  $(x \ominus y) \ominus ((y \ominus z) \ominus (x \ominus z)) = 1$ ,
- (a3)  $x \ominus x = 1$ ,
- (a4)  $x \ominus y = y \ominus x = 1 \Rightarrow x = y$ .

We call such algebra a *WFI-algebra*. A nonempty subset  $S$  of  $G$  is called a *subalgebra* of  $\mathcal{G}$  if  $x \ominus y \in S$  whenever  $x, y \in S$ . A nonempty subset  $F$  of  $G$  is called a *filter* of  $\mathcal{G}$  if it satisfies:

- $1 \in F$ ,
- $x \ominus y \in F$  and  $x \in F$  imply  $y \in F$  for all  $x, y \in G$ .

In a WFI-algebra  $\mathcal{G}$ , the following are true (see [1]):

- (b1)  $x \ominus ((x \ominus y) \ominus y) = 1$ ,
- (b2)  $1 \ominus x = 1 \Rightarrow x = 1$ ,
- (b3)  $1 \ominus x = x$ ,
- (b4)  $x \ominus y = 1 \Rightarrow (y \ominus z) \ominus (x \ominus z) = 1, (z \ominus x) \ominus (z \ominus y) = 1$ ,
- (b5)  $(x \ominus y) \ominus 1 = (x \ominus 1) \ominus (y \ominus 1)$ .

We now define a relation " $\preceq$ " on  $\mathcal{G}$  by  $x \preceq y$  if and only if  $x \ominus y = 1$ . It is easy to verify that a WFI-algebra is a partially ordered set with respect to  $\preceq$ . A WFI-algebra  $\mathcal{G}$  is said to be *associative* [1] if it satisfies

$(x \ominus y) \ominus z = x \ominus (y \ominus z)$  for all  $x, y, z \in G$ . A WFI-algebra  $\mathcal{G}$  is said to be *medial* [1] if it satisfies

$$(x \ominus y) \ominus (a \ominus b) = (x \ominus a) \ominus (y \ominus b)$$

for all  $x, y, a, b \in G$ .

### 3. Simulative and Mutant WFI-algebras

For a WFI-algebra  $\mathcal{G}$ , the set

$$\mathcal{S}(\mathcal{G}) := \{x \in G \mid x \preceq 1\}$$

is called the *simulative part* of  $\mathcal{G}$ . Let  $x, y \in \mathcal{S}(\mathcal{G})$ . Then  $x \ominus 1 = 1$  and  $y \ominus 1 = 1$ . It follows from (b5) that

$$(x \ominus y) \ominus 1 = (x \ominus 1) \ominus (y \ominus 1) = 1 \ominus 1 = 1$$

so that  $x \ominus y \preceq 1$ . Hence  $x \ominus y \in \mathcal{S}(\mathcal{G})$ , which shows that  $\mathcal{S}(\mathcal{G})$  is a subalgebra of  $\mathcal{G}$ .

**Proposition 3.1.** *The simulative part of a WFI-algebra  $\mathcal{G}$  is a filter of  $\mathcal{G}$ .*

*Proof.* Obviously,  $1 \in \mathcal{S}(\mathcal{G})$ . Let  $x, y \in G$  be such that  $x \in \mathcal{S}(\mathcal{G})$  and  $x \ominus y \in \mathcal{S}(\mathcal{G})$ . Then  $x \preceq 1$  and  $x \ominus y \preceq 1$ . It follows from (b5) and (b3) that

$$1 = (x \ominus y) \ominus 1 = (x \ominus 1) \ominus (y \ominus 1) = 1 \ominus (y \ominus 1) = y \ominus 1,$$

that is,  $y \preceq 1$ . Hence  $y \in \mathcal{S}(\mathcal{G})$ , and the proof is complete. □

**Definition 3.2.** A WFI-algebra  $\mathcal{G}$  is said to be *simulative* if it satisfies

$$(S) \quad x \preceq 1 \Rightarrow x = 1.$$

Note that the condition (S) is equivalent to  $\mathcal{S}(\mathcal{G}) = \{1\}$ .

**Example 3.3.** (1) Let  $\mathbb{Z}$  be the set of integers. Then  $\mathfrak{Z} := (\mathbb{Z}; \ominus, 0)$  is a simulative WFI-algebra, where  $x \ominus y = y - x$  for all  $x, y \in \mathbb{Z}$ .

(2) Let  $G = \{1, a, b\}$  be a set with the following Cayley table and Hasse diagram:

$\ominus$	1	a	b
1	1	a	b
a	1	1	b
b	b	b	1

Then  $\mathcal{G} := (G; \ominus, 1)$  is a WFI-algebra (see [1]) which does not satisfy the condition (S).

**Theorem 3.4.** *Let  $\mathcal{G}$  be a WFI-algebra. Then the following are equivalent.*

- (i)  $\mathcal{G}$  is simulative.
- (ii)  $(x \ominus 1) \ominus 1 = x, \forall x \in G$ .
- (iii)  $(x \ominus 1) \ominus y = (y \ominus 1) \ominus x, \forall x, y \in G$ .

*Proof.* (i)  $\Rightarrow$  (ii) Suppose that  $\mathcal{G}$  is simulative. Using (a1) and (a3), we have

$$x \ominus ((x \ominus 1) \ominus 1) = (x \ominus 1) \ominus (x \ominus 1) = 1,$$

that is,  $x \preceq (x \ominus 1) \ominus 1$ . It follows from (a3) and (b4) that

$$((x \ominus 1) \ominus 1) \ominus x \preceq x \ominus x = 1$$

so from (S) that  $((x \ominus 1) \ominus 1) \ominus x = 1$ , that is,  $(x \ominus 1) \ominus 1 \preceq x$ . Hence, by (a4), we get  $(x \ominus 1) \ominus 1 = x$ .

(ii)  $\Rightarrow$  (iii) Assume that  $\mathcal{G}$  satisfies (ii). Then

$$\begin{aligned} (x \ominus 1) \ominus y &= (x \ominus 1) \ominus ((y \ominus 1) \ominus 1) && \text{by (ii)} \\ &= (y \ominus 1) \ominus ((x \ominus 1) \ominus 1) && \text{by (a1)} \\ &= (y \ominus 1) \ominus x, && \text{by (ii)} \end{aligned}$$

which proves (iii).

(iii)  $\Rightarrow$  (i) Suppose that  $\mathcal{G}$  satisfies (iii) and let  $x \in G$  be such that  $x \preceq 1$ . Then

$$x = 1 \ominus x = (1 \ominus 1) \ominus x = (x \ominus 1) \ominus 1 = 1 \ominus 1 = 1,$$

and so  $\mathcal{G}$  is simulative.  $\square$

**Theorem 3.5.** *Let  $\mathcal{G}$  be a WFI-algebra. Then the following are equivalent.*

- (i)  $\mathcal{G}$  is simulative.
- (ii)  $x \preceq y \Rightarrow x = y$ .
- (iii)  $(x \ominus y) \ominus (z \ominus y) = z \ominus x, \forall x, y, z \in G$ .
- (iv)  $(x \ominus y) \ominus 1 = y \ominus x, \forall x, y \in G$ .
- (v)  $\mathcal{G}$  satisfies the right cancellation law, i.e.,  $x \ominus z = y \ominus z \Rightarrow x = y$ .
- (vi)  $(x \ominus y) \ominus (z \ominus y) = (x \ominus z) \ominus 1, \forall x, y, z \in G$ .

*Proof.* (i)  $\Rightarrow$  (ii) Suppose that  $\mathcal{G}$  is simulative and let  $x, y \in G$  be such that  $x \preceq y$ . Then

$$\begin{aligned} (y \ominus x) \ominus 1 &= (y \ominus 1) \ominus (x \ominus 1) && \text{by (b5)} \\ &= (y \ominus (x \ominus y)) \ominus (x \ominus 1) && \text{since } x \preceq y \\ &= (x \ominus (y \ominus y)) \ominus (x \ominus 1) && \text{by (a1)} \\ &= (x \ominus 1) \ominus (x \ominus 1) && \text{by (a3)} \\ &= 1, && \text{by (a3)} \end{aligned}$$

and so  $y \ominus x \in \mathcal{S}(\mathcal{G}) = \{1\}$ . Thus  $y \ominus x = 1$ , i.e.,  $y \preceq x$ . It follows from (a4) that  $x = y$ .

(ii)  $\Rightarrow$  (iii) Assume that (ii) holds. Since  $z \ominus x \preceq (x \ominus y) \ominus (z \ominus y)$ , it follows from (ii) that  $z \ominus x = (x \ominus y) \ominus (z \ominus y)$ .

(iii)  $\Rightarrow$  (iv) Suppose that (iii) is true. Then

$$(x \ominus y) \ominus 1 = (x \ominus 1) \ominus (y \ominus 1) = y \ominus x.$$

(iv)  $\Rightarrow$  (v) Assume that (iv) is true and let  $x, y, z \in G$  be such that  $x \ominus z = y \ominus z$ . Then

$$y \ominus x = (x \ominus y) \ominus 1 = (x \ominus y) \ominus ((y \ominus z) \ominus (x \ominus z)) = 1.$$

Similarly, we have  $x \ominus y = 1$ , and so  $x = y$  by (a4).

(v)  $\Rightarrow$  (i) Suppose that  $\mathcal{G}$  satisfies the right cancellation law. Let  $x \in \mathcal{S}(\mathcal{G})$ . Then  $x \ominus 1 = 1 = 1 \ominus 1$ , and so  $x = 1$  by using (v). Therefore  $\mathcal{S}(\mathcal{G}) = \{1\}$  which shows that  $\mathcal{G}$  is simulative.

(iv)  $\Rightarrow$  (vi) If (iv) is true, then (iii) is valid, and so

$$(x \ominus y) \ominus (z \ominus y) = z \ominus x = (x \ominus z) \ominus 1.$$

(vi)  $\Rightarrow$  (i) Suppose that (vi) is valid and let  $x \in \mathcal{S}(\mathcal{G})$ . Then  $x \ominus 1 = 1 = x \ominus x$ , which implies that

$$x = (x \ominus x) \ominus (1 \ominus x) = (x \ominus 1) \ominus 1 = 1 \ominus 1 = 1.$$

This shows that  $\mathcal{S}(\mathcal{G}) = \{1\}$ . Hence  $\mathcal{G}$  is simulative. This completes the proof.  $\square$

**Proposition 3.6.** *Let  $\mathcal{G}$  be a simulative WFI-algebra. Then*

$$x \ominus (y \ominus z) = ((x \ominus 1) \ominus y) \ominus z, \forall x, y, z \in G.$$

*Proof.* For any  $x, y, z \in G$ , we have

$$\begin{aligned} x \ominus (y \ominus z) &= ((x \ominus 1) \ominus 1) \ominus (y \ominus z) && \text{by Theorem 3.4} \\ &= y \ominus (((x \ominus 1) \ominus 1) \ominus z) && \text{by (a1)} \\ &= y \ominus ((z \ominus 1) \ominus (x \ominus 1)) && \text{by Theorem 3.4} \\ &= (z \ominus 1) \ominus (y \ominus (x \ominus 1)) && \text{by (a1)} \\ &= ((y \ominus (x \ominus 1)) \ominus 1) \ominus z && \text{by Theorem 3.4} \\ &= ((x \ominus 1) \ominus y) \ominus z, && \text{by Theorem 3.5} \end{aligned}$$

which completes the proof.  $\square$

**Theorem 3.7.** *Let  $\mathcal{G}$  be a simulative WFI-algebra and define a binary operation “ $\cdot$ ” on  $\mathcal{G}$  by  $x \cdot y = (y \ominus 1) \ominus x$  for all  $x, y \in G$ . Then  $(G, \cdot, 1)$  is a commutative group.*

*Proof.* Using (a1) and Theorem 3.4(iii), we have

$$\begin{aligned} x \cdot (y \cdot z) &= x \cdot ((z \ominus 1) \ominus y) = (((z \ominus 1) \ominus y) \ominus 1) \ominus x \\ &= (x \ominus 1) \ominus ((z \ominus 1) \ominus y) = (z \ominus 1) \ominus ((x \ominus 1) \ominus y) \\ &= (z \ominus 1) \ominus ((y \ominus 1) \ominus x) = (x \cdot y) \cdot z, \end{aligned}$$

and  $x \cdot y = (y \ominus 1) \ominus x = (x \ominus 1) \ominus y = y \cdot x$ . It follows from Theorem 3.4(ii) that  $x \cdot 1 = 1 \cdot x = (x \ominus 1) \ominus 1 = x$  and

$$(x \ominus 1) \cdot x = x \cdot (x \ominus 1) = ((x \ominus 1) \ominus 1) \ominus x = x \ominus x = 1.$$

Hence  $(G, \cdot, 1)$  is a commutative group with  $x \ominus 1$  as the inverse of  $x$ .  $\square$

Conversely, we have the following.

**Theorem 3.8.** *If  $(G, \cdot, 1)$  is a commutative group, then  $\mathcal{G} := (G; \ominus, 1)$  is a simulative WFI-algebra, where  $x \ominus y = x^{-1} \cdot y$  for all  $x, y \in G$ .*

*Proof.* It is easy to verify the axioms of a WFI.  $\square$

**Lemma 3.9.** [1, Proposition 3.25] *Every medial WFI-algebra satisfies the following identities:*

- (i)  $(x \ominus y) \ominus 1 = y \ominus x$ .
- (ii)  $(x \ominus 1) \ominus 1 = x$ .
- (iii)  $(x \ominus y) \ominus y = x$ .

**Theorem 3.10.** *Let  $\mathcal{G}$  be a WFI-algebra. Then the following are equivalent:*

- (i)  $\mathcal{G}$  is simulative.
- (ii)  $(x \ominus y) \ominus y = x, \forall x, y \in G$ .
- (iii)  $\mathcal{G}$  is medial.

*Proof.* (iii)  $\Rightarrow$  (ii) and (ii)  $\Rightarrow$  (i) are by Lemma 3.9 and Theorem 3.4.

(i)  $\Rightarrow$  (iii) Assume that  $\mathcal{G}$  is simulative. Note that, for all  $x, y \in G$ ,

$$(x \cdot y) \ominus 1 = (x \ominus 1) \cdot (y \ominus 1) = x \ominus (y \ominus 1)$$

and  $x \ominus y = ((x \ominus 1) \ominus 1) \ominus y = (x \ominus 1) \cdot y$ . It follows that

$$\begin{aligned}
 (x \ominus y) \ominus (a \ominus b) &= ((x \ominus 1) \cdot y) \ominus ((a \ominus 1) \cdot b) \\
 &= (((x \ominus 1) \cdot y) \ominus 1) \cdot ((a \ominus 1) \cdot b) \\
 &= ((x \ominus 1) \ominus (y \ominus 1)) \cdot ((a \ominus 1) \cdot b) \\
 &= (((x \ominus 1) \ominus 1) \cdot (y \ominus 1)) \cdot ((a \ominus 1) \cdot b) \\
 &= x \cdot (y \ominus 1) \cdot (a \ominus 1) \cdot b \\
 &= x \cdot (a \ominus 1) \cdot (y \ominus 1) \cdot b \\
 &= (((x \ominus 1) \ominus 1) \cdot (a \ominus 1)) \cdot ((y \ominus 1) \cdot b) \\
 &= ((x \ominus 1) \ominus (a \ominus 1)) \cdot ((y \ominus 1) \cdot b) \\
 &= (((x \ominus 1) \cdot a) \ominus 1) \cdot ((y \ominus 1) \cdot b) \\
 &= ((x \ominus 1) \cdot a) \ominus ((y \ominus 1) \cdot b) \\
 &= (x \ominus a) \ominus (y \ominus b)
 \end{aligned}$$

for all  $x, y, a, b \in G$ . Hence  $\mathcal{G}$  is medial.  $\square$

**Theorem 3.11.** *Every associative WFI-algebra is a simulative WFI-algebra.*

*Proof.* Let  $\mathcal{G}$  be an associative WFI-algebra. It suffices to show that  $\mathcal{G}$  satisfies the identity  $(x \ominus y) \ominus y = x$  for all  $x, y \in G$  (see Theorem 3.10). Obviously  $x \preceq (x \ominus y) \ominus y$ . Now, using (a1), (a3), and the associativity, we have

$$\begin{aligned}
 ((x \ominus y) \ominus y) \ominus x &= (x \ominus y) \ominus (y \ominus x) = y \ominus ((x \ominus y) \ominus x) \\
 &= y \ominus (x \ominus (y \ominus x)) = (y \ominus x) \ominus (y \ominus x) = 1,
 \end{aligned}$$

that is,  $(x \ominus y) \ominus y \preceq x$ . It follows from (a4) that  $(x \ominus y) \ominus y = x$ . Hence  $\mathcal{G}$  is simulative.  $\square$

The converse of Theorem 3.11 may not be true as shown in the following example.

**Example 3.12.** Let  $G = \{1, a, b\}$  be a set with the following Cayley table and Hasse diagram:



$\ominus$	1	a	b	
1	1	a	b	
a	b	1	a	$\begin{matrix} \circ & \circ & \circ \\ 1 & a & b \end{matrix}$
b	a	b	1	

Then  $\mathcal{G} := (G; \ominus, 1)$  is a simulative WFI-algebra, but it is not an associative WFI-algebra since  $b \ominus (a \ominus 1) \neq (b \ominus a) \ominus 1$ .

For a WFI-algebra  $\mathcal{G}$ , consider the set

$$\mathcal{L}(\mathcal{G}) := \{x \in G \mid x \ominus 1 = x\}.$$

Note that  $\mathcal{L}(\mathcal{G})$  is a subalgebra of  $\mathcal{G}$  (see [1]).

**Theorem 3.13.** *Let  $\mathcal{G}$  be a WFI-algebra such that  $\mathcal{L}(\mathcal{G}) = G$ . Then the following are equivalent:*

- (i)  $\mathcal{G}$  is associative.
- (ii)  $x \ominus (x \ominus y) = y, \forall x, y \in G$ .
- (iii)  $x \ominus (y \ominus z) = z \ominus (y \ominus x), \forall x, y, z \in G$ .
- (iv)  $(x \ominus y) \ominus y = x, \forall x, y \in G$ .
- (v)  $\mathcal{G}$  is medial.
- (vi)  $\mathcal{G}$  is simulative.

*Proof.* (i)  $\Rightarrow$  (ii) is by (a3), (b3) and the associativity.

(ii)  $\Rightarrow$  (iii) Note that

$$\begin{aligned} 1 &= x \ominus x && \text{by (a3)} \\ &= x \ominus (z \ominus (z \ominus x)) && \text{by (ii)} \\ &= x \ominus (z \ominus ((z \ominus x) \ominus 1)) && \text{since } \mathcal{L}(\mathcal{G}) = G \\ &= x \ominus ((z \ominus x) \ominus (z \ominus 1)) && \text{by (a1)} \\ &= x \ominus ((z \ominus x) \ominus z) && \text{since } \mathcal{L}(\mathcal{G}) = G \\ &\preceq x \ominus ((y \ominus (z \ominus x)) \ominus (y \ominus z)) && \text{by (a1), (a2) and (b4)} \\ &= x \ominus ((z \ominus (y \ominus x)) \ominus (y \ominus z)) && \text{by (a1)} \\ &= (z \ominus (y \ominus x)) \ominus (x \ominus (y \ominus z)). && \text{by (a1)} \end{aligned}$$

Hence  $z \ominus (y \ominus x) \preceq x \ominus (y \ominus z)$ . Similarly, we have  $x \ominus (y \ominus z) \preceq z \ominus (y \ominus x)$  by symmetry. It follows from (a4) that  $x \ominus (y \ominus z) = z \ominus (y \ominus x)$ .

(iii)  $\Rightarrow$  (v) Suppose that (iii) holds. Then

$$\begin{aligned}(x \ominus y) \ominus (a \ominus b) &= b \ominus (a \ominus (x \ominus y)) = b \ominus (y \ominus (x \ominus a)) \\ &= (x \ominus a) \ominus (y \ominus b),\end{aligned}$$

and so  $\mathcal{G}$  is medial.

(iv)  $\Leftrightarrow$  (v)  $\Leftrightarrow$  (vi) See Theorem 3.10.

(iv)  $\Rightarrow$  (1) Assume that  $\mathcal{G}$  is simulative. Then  $(G, \cdot, 1)$  is a commutative group (see Theorem 3.7). Using the condition  $\mathcal{L}(\mathcal{G}) = G$ , we have

$$x \ominus y = (x \ominus 1) \ominus y = y \cdot x = x \cdot y = (y \ominus 1) \ominus x = y \ominus x,$$

and hence  $x \ominus (y \ominus z) = x \ominus (z \ominus y) = z \ominus (x \ominus y) = (x \ominus y) \ominus z$ . Therefore  $\mathcal{G}$  is associative. This completes the proof.  $\square$

Let  $\mathcal{G}$  be a WFI-algebra. For nonnegative integers  $i$  and  $j$ , we define a polynomial  $P_{i,j}(x, y)$  of two variables  $x$  and  $y$  in  $G$  as follows:

$$\begin{aligned}P_{0,0}(x, y) &= (y \ominus x) \ominus x, \\ P_{i+1,j}(x, y) &= (y \ominus x) \ominus P_{i,j}(x, y), \\ P_{i,j+1}(x, y) &= (x \ominus y) \ominus P_{i,j}(x, y).\end{aligned}$$

**Definition 3.14.** Let  $i, j, m$  and  $n$  be nonnegative integers. A WFI-algebra  $\mathcal{G}$  is said to be  $(i, j; m, n)$ -mutant if  $P_{i,j}(x, y) = P_{m,n}(y, x)$  for all  $x, y \in G$ .

In particular, if  $i = j = m = n = 0$ , then we say that  $\mathcal{G}$  is mutant, that is, a  $(0, 0; 0, 0)$ -mutant WFI-algebra is called a mutant WFI.

**Example 3.15.** The WFI-algebra  $\mathcal{G}$  in Example 3.3(2) is a  $(1, 0; 0, 0)$ -mutant WFI-algebra, but the WFI-algebra  $\mathfrak{Z} := (\mathbb{Z}; \ominus, 0)$  in Example 3.3(1) may not be  $(1, 0; 0, 0)$ -mutant because  $P_{1,0}(2, 3) = 4 \neq 2 = P_{0,0}(3, 2)$ .

**Theorem 3.16.** Every simulative WFI-algebra  $\mathcal{G}$  is  $(0, n + 1; n, 0)$ -mutant for every nonnegative integer  $n$ .

*Proof.* The proof is by induction on  $n$ . Let  $x, y \in G$ . If  $n = 0$ , then

$$\begin{aligned} P_{0,1}(x, y) &= (x \ominus y) \ominus P_{0,0}(x, y) \\ &= (x \ominus y) \ominus ((y \ominus x) \ominus x) \\ &= (x \ominus y) \ominus y \\ &= P_{0,0}(y, x). \end{aligned}$$

Assume that the result is valid for  $n = k$ , i.e.,  $P_{0,k+1}(x, y) = P_{k,0}(y, x)$ .

Then

$$\begin{aligned} P_{0,k+2}(x, y) &= (x \ominus y) \ominus P_{0,k+1}(x, y) \\ &= (x \ominus y) \ominus P_{k,0}(y, x) \\ &= P_{k+1,0}(y, x). \end{aligned}$$

Hence  $P_{0,n+1}(x, y) = P_{n,0}(y, x)$ , that is,  $\mathcal{G}$  is  $(0, n + 1; n, 0)$ -mutant.  $\square$

The second part of Example 3.15 shows that a simulative WFI-algebra  $\mathcal{G}$  may not be  $(1, 0; 0, 0)$ -mutant.

**Theorem 3.17.** *Every simulative WFI-algebra is  $(\alpha, m; m, \alpha + 1)$ -mutant for all nonnegative integer  $m$ , where  $\alpha = m + i$  for  $i \geq 0$ .*

*Proof.* Let  $\mathcal{G}$  be a simulative WFI-algebra. We first show that for every nonnegative integer  $n$ ,

- (i)  $P_{n,n}(x, y) = y$ ,
- (ii)  $P_{n,n+1}(x, y) = x$ ,
- (iii)  $P_{n+1,n}(x, y) = (y \ominus x) \ominus y$ ,
- (iv) If we use the notation  $x^k \ominus y$  instead of

$$\overbrace{x \ominus (\cdots \ominus (x \ominus (x \ominus y)) \cdots)}^{k\text{-times}},$$

then  $P_{n+i,n}(x, y) = (y \ominus x)^i \ominus y = P_{n,n+i+1}(y, x)$  for  $i \geq 0$ .

By induction on  $n$ , we have

$$P_{0,0}(x, y) = (y \ominus x) \ominus x = y$$

by Theorem 3.10. Suppose that  $P_{k,k}(x, y) = y$  for  $n = k > 0$ . Then

$$\begin{aligned} P_{k+1,k+1}(x, y) &= (y \ominus x) \ominus P_{k,k+1}(x, y) \\ &= (y \ominus x) \ominus ((x \ominus y) \ominus P_{k,k}(x, y)) \\ &= (y \ominus x) \ominus ((x \ominus y) \ominus y) \\ &= (y \ominus x) \ominus x = y, \end{aligned}$$

which proves (i). Note that

$$\begin{aligned} P_{0,1}(x, y) &= (x \ominus y) \ominus P_{0,0}(x, y) \\ &= (x \ominus y) \ominus ((y \ominus x) \ominus x) \\ &= (x \ominus y) \ominus y = x. \end{aligned}$$

Suppose that (ii) is valid for  $n = k > 0$ . Then

$$\begin{aligned} P_{k+1,k+2}(x, y) &= (x \ominus y) \ominus P_{k+1,k+1}(x, y) \\ &= (x \ominus y) \ominus ((y \ominus x) \ominus P_{k,k+1}(x, y)) \\ &= (x \ominus y) \ominus ((y \ominus x) \ominus x) \\ &= (x \ominus y) \ominus y = x. \end{aligned}$$

Hence (ii) holds. We observe that

$$P_{1,0}(x, y) = (y \ominus x) \ominus P_{0,0}(x, y) = (y \ominus x) \ominus y.$$

If (iii) is true for  $n = k > 0$ , then

$$P_{k+2,k+1}(x, y) = (y \ominus x) \ominus P_{k+1,k+1}(x, y) = (y \ominus x) \ominus y$$

by (i). Therefore (iii) is valid. For a fixed  $n \geq 0$ , we have

$$P_{n+0,n}(x, y) = P_{n,n}(x, y) = y = (y \ominus x)^0 \ominus y.$$

Assume that  $P_{n+k,n}(x, y) = (y \ominus x)^k \ominus y$  for  $i = k > 0$ . Then

$$\begin{aligned} P_{n+k+1,n}(x, y) &= (y \ominus x) \ominus P_{n+k,n}(x, y) \\ &= (y \ominus x) \ominus ((y \ominus x)^k \ominus y) \\ &= (y \ominus x)^{k+1} \ominus y. \end{aligned}$$

From (ii), it follows that

$$P_{n,n+1}(y, x) = y = (y \ominus x)^0 \ominus y.$$

Suppose that  $P_{n,n+k+1}(y, x) = (y \ominus x)^k \ominus y$  for  $i = k > 0$ . Then

$$\begin{aligned} P_{n,n+k+2}(y, x) &= (y \ominus x) \ominus P_{n,n+k+1}(y, x) \\ &= (y \ominus x) \ominus ((y \ominus x)^k \ominus y) \\ &= (y \ominus x)^{k+1} \ominus y. \end{aligned}$$

This proves (iv). Finally for a fixed  $m \geq 0$ , let  $\alpha = m + i$  where  $i \geq 0$ . Using (iv), we conclude that  $P_{\alpha,m}(x, y) = P_{m,\alpha+1}(y, x)$ . Therefore  $\mathcal{G}$  is  $(\alpha, m; m, \alpha + 1)$ -mutant. This completes the proof.  $\square$

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#### References

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