# A New Adaptive Load Sharing Mechanism in Homogeneous Distributed Systems Using Genetic Algorithm

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#### **ABSTRACT**

Load sharing is a critical resource in computer system. In sender-initiated load sharing algorithms, the sender continues to send unnecessary request messages for load transfer until a receiver is found while the system load is heavy. Meanwhile, in the receiver-initiated load sharing algorithms, the receiver continues to send an unnecessary request message for load acquisition until a sender is found while the system load is light. These unnecessary request messages result in inefficient communications, low CPU utilization, and low system throughput in distributed systems. To solve these problems, we propose a genetic algorithm based approach for improved sender-initiated and receiver-initiated load sharing in distributed systems. And we expand this algorithm to an adaptive load sharing algorithm. Compared with the conventional sender-initiated and receiver-initiated algorithms, the proposed algorithm decreases the response time and task processing time.

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#### 1. INTRODUCTION

Distributed systems consist of a collection of autonomous computers connected network. The primary advantages of these systems are high performance, availability, and extensibility at low cost. To improve a performance of distributed systems, it is essential to keep the system load to each processor equally.

An objective of load sharing in distributed systems is to allocate tasks among the processors to maximize the utilization of processors and to minimize the mean response time. Load sharing algorithms can be largely classified into three classes: static, dynamic, adaptive. Our approach is based on dynamic load sharing algorithm. In dynamic scheme, an overloaded processor(sender) sends excess tasks to an underloaded processor(receiver) during execution.

Dynamic load sharing algorithms are specialized into three methods: sender-initiated, receiver-initiated, symmetrically-initiated. Basically our approach is a sender-initiated and receiver-initiated algorithm.

Under sender-initiated algorithms, load sharing activity is initiated by a sender trying to send a task to a receiver [1,2]. Decision of task transfer is made in each processor independently. A request message for the task transfer is initially issued from a sender to an another processor randomly selected. If the selected processor is receiver, it returns an accept message. And the receiver is ready to receive an additional task from sender. Otherwise, it returns a reject message, and the sender tries for others until receiving an

Similarly, Under receiver-initiated algorithms, load sharing activity is initiated by a receiver trying to receive a task from a sender [1,2]. Decision of task acquisition is made in each processor independently. A request message for the task acquisition is initially issued from a receiver to an another processor randomly selected. If the selected processor is sender. it returns an accept message. And the sender is ready to transfer an additional task to receiver. Otherwise, it returns a reject message, and the receiver tries for others until receiving an accept message. While distributed systems remain to high system load, the receiver-initiated algorithm performs well. But when the distributed systems become to light system load, it is difficult to find a suitable sender because most processors have small tasks. So, many request and reject messages are repeatedly sent back and forth, and a lot of time is consumed before execution. Therefore, much of the task processing time is consumed, and causes low system throughput, low CPU utilization

To solve these problems in sender-initiated and receiver-initiated algorithm, we use a new genetic algorithm and expand to a new adaptive load sharing algorithm. A new genetic algorithm evolves strategy for determining a destination processor to receive a task in sender-initiated algorithm and to send a task in receiver-initiated algorithm. And we define a

accept message. While distributed systems remain to light system load, the sender-initiated algorithm performs well. But when the distributed systems become to heavy system load, it is difficult to find a suitable receiver because most processors have additional tasks to send. So, many request and reject messages are repeatedly sent back and forth, and a lot of time is consumed before execution.

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suitable fitness function. In this scheme, a number of request messages issued before accepting a task are determined through a proposed genetic algorithm. The proposed genetic algorithm applies to a population of binary strings. Each gene in the string stands for a number of processors which request messages should be sent off.

#### 2. GA-BASED METHOD

In this section, we describe various factors to be needed for GA-based load sharing. These are load measure, coding method, fitness function and algorithm.

#### 2.1 Load Measure

We employ the CPU queue length as a suitable load index because this measure is known as a most suitable index[5]. This measure means a number of tasks in CPU queue residing in a processor.

We use a 3-level scheme to represent a load state on its own CPU queue length of a processor. Table 1 shows the 3-level load measurement scheme.  $T_{up}$  and  $T_{low}$  are algorithm design parameters and are called upper and lower thresholds respectively

Table 1. 3-level load measurement scheme

Load state	Meaning	Criteria
L-load	light-load	$CQL \leq T_{low}$
N-load	normal-load	$T_{low} < \text{CQL} \le T_{up}$
H-load	heavy-load	$CQL > T_{up}$

( CQL : CPU Queue Length )

Transfer policy uses the threshold policy that makes decision based on CPU queue length. The transfer policy is triggered when a task arrives. A node identifies as a sender if a new task originating at the node makes the CPU queue length exceed  $T_{up}$ . A node identifies itself as a suitable receiver for a task acquisition if the node's CPU queue length will not cause to exceed  $T_{low}$ .

Each processor in distributed systems has its own population which genetic operators are applied to. There are many encoding methods; Binary encoding, Character and real-valued encoding, Tree encoding. We use binary encoding method in this paper. So, a string in population can be defined as a binary-coded vector  $\langle v_o, v_l, ..., v_{n-l} \rangle$  which indicates a set of processors to which the request messages are sent off. If the request message is transferred to the processor  $P_i$ (where  $0 \le i \le n-1$ , n is the total number of processors), then  $v_i = 1$ , otherwise  $v_i = 0$ . Each string has its own fitness value. We select a string by a probability proportional to its fitness value, and transfer the request messages to the processors indicated by the string.

## 2.2 Sender-based Load Sharing Approach

In the sender-based load sharing approach using genetic algorithm, processors received the request message from the sender send accept message or reject message depending on its own CPU queue length. In the case of more than two accept

messages returned, one is selected at random.

Suppose that there are 10 processors in distributed systems, and the processor  $P_0$  is a sender. Then, genetic algorithm is performed to decide a suitable receiver. It is selected a string by a probability proportional to its fitness value. Suppose a selected string is <-, 1, 0, 1, 0, 0, 1, 1, 0, 0>, then the sender  $P_0$  sends request messages to the processors  $(P_1, P_3, P_6, P_7)$ . After each processor  $(P_1, P_3, P_6, P_7)$  receives a request message from the processor  $P_0$ , each processor checks its load state. If the processor  $P_3$  is a light load state, the processor  $P_3$  sends back an accept message to the processor  $P_0$ . Then the processor  $P_0$  transfers a task to the processor  $P_3$ .

Each string included in a population is evaluated by the fitness function using following formula.

```
F_i = 1 / ((\alpha \times TMP) + (\beta \times TMT) + (\gamma \times TTP))
```

Here,  $\alpha$ ,  $\beta$ ,  $\gamma$  mean the weights for parameters such as *TMP*, *TMT*, *TTP*. The purpose of the weights is to be operated equally for each parameter to fitness function  $F_i$ .

TMP(Total Message Processing time) is the summation of the processing times for request messages to be transferred. TMT(Total Message Transfer time) means the summation of each message transfer times from the sender to processors corresponding to bits set '1' in selected string. The objective of this parameter is to select a string with the shortest distance eventually. TTP(Total Task Processing time) is the summation of the times needed to perform a task at each processor corresponding to bits set '1' in selected string. The objective of this parameter is to select a string with the fewest loads. Eventually, a string with the largest fitness value in population is selected. And after genetic\_operation is performed, the request messages are transferred to processors corresponding to bits set '1' in selected string.

This algorithm consists of five modules such as Initialization, Check\_load, String\_evaluation, Genetic\_operation and Message\_evaluation. Genetic\_operation module consists of three sub-modules such as Local improvement, Reproduction, Crossover. These modules are executed at each processor in distributed systems. The algorithm of the proposed sender-based load sharing is as following.

```
/* GA-based sender-initiated load sharing algorithm */
Procedure Genetic_algorithm Approach

{ Initialization()
    while (Check_load())
    if (Loadi > Tup) {
        Individual_evaluation();
        Genetic_operation();
        Message_evaluation();
    }

Process a task in local processor;
}

Procedure Genetic_operation()
{ Local_improvement_operation();
    Reproduction(); Crossover(); }
```

An Initialization module is executed in each processor. A

population of strings is randomly generated without duplication. A Check\_load module is used to observe its own processor's load by checking the CPU queue length, whenever a task is arrived in a processor. If the observed load is heavy, the load sharing algorithm performs the above modules. A String\_evaluation module calculates the fitness value of strings in the population. A Genetic\_operation module such as Local improvement, Reproduction, Crossover is executed on the population in such a way as above. Distributed systems consist of groups with autonomous computers. When each group consists of many processors, we can suppose that there are p parts in a string corresponding to the groups. The following genetic operations are applied to each strings, and new population of strings is generated:

## (1) Local Improvement Operation

String 1 is chosen. A copy version of the string 1 is generated and part 1 of the newly generated string is mutated. This new string is evaluated by proposed fitness function. If the evaluated value of the new string is higher than that of the original string, replace the original string with the new string. After this, the local improvement of part 2 of string 1 is done repeatedly. This local improvement is applied to each part one by one. When the local improvement of all the parts is finished, new string 1 is generated. String 2 is then chosen, and the above-mentioned local improvement is done. This local improvement operation is applied to all the strings in population.

/\* Algorithms for local improvement operation \*/

## (2) Reproduction

The reproduction operation is applied to the newly generated strings. We use the "wheel of fortune" technique[4].

## (3) Crossover

The crossover operation is applied to the newly generated strings. These newly generated strings are evaluated. We applied to the "one-point" crossover operator in this paper [4]. One-point crossover used in this paper differs from the pure one-point crossover operator. In pure one-point crossover, crossover activity generates based on randomly selected crossover point in the string. But boundaries between parts(p) are used as an alternative of crossover points in this paper. So we select a boundary among many boundaries at random. And

a selected boundary is used as a crossover point. This purpose is to preserve effect of the local improvement operation of the previous phase.

Suppose that there are 5 parts in distributed systems. A boundary among the many boundaries  $(B_1, B_2, B_3, B_4)$  is determined at random as a crossover point. If a boundary  $B_3$  is selected as a crossover point, crossover activity generate based on the  $B_3$ . So, the effect of the local improvement operation in the previous phase is preserved through crossover activity.

The Genetic\_operation selects a string from the population at the probability proportional to its fitness, and then sends off the request messages according to the contents of the selected string.

A Message\_evaluation module is used whenever a processor receives a message from other processors. When a processor  $P_i$  receives a request message, it sends back an accept or reject message depending on its CPU queue length.

#### 2.3 Receiver-based Load Sharing Approach

In the receiver-initiated load sharing approach, a receiver finds a sender to obtain a task. Processors received the request message from the receiver send accept message or reject message depending on its own CPU queue length. In the case of more than two accept messages returned, one is selected at random.

Suppose that there are 10 processors in distributed systems, and the processor  $P_0$  is a receiver. Then, genetic algorithm is performed to decide a suitable sender. It is selected a string by a probability proportional to its fitness value. Suppose a selected string is <-, 1, 0, 1, 0, 0, 1, 1, 0, 0>, then the receiver  $P_0$  sends request messages to the processors ( $P_1$ ,  $P_3$ ,  $P_6$ ,  $P_7$ ). After each processor( $P_1$ ,  $P_3$ ,  $P_6$ ,  $P_7$ ) receives a request message from the processor  $P_0$ , each processor checks its load state. If the processor  $P_3$  is a heavy load state, the processor  $P_3$  sends back an accept message to the processor  $P_0$ . Then the processor  $P_0$  receives a task from the processor  $P_3$ .

Each string included in a population is evaluated by the fitness function using following formula in the receiver-initiated approach.

$$F_i = 1 / ((\alpha \times TMP) + (\beta \times TMT)) + (\gamma \times TTP))$$

So, in order to have a largest fitness value, each parameter such as *TMP*, *TMT* must have small values as possible as, but *TTP* must have large value. That is, *TMP* must have fewer number for request messages, and *TMT* must have the shortest distance, and *TTP* should have the larger number of tasks. Eventually, in GA-based receiver-initiated load sharing algorithm, a receiver easily must find a sender. Therefore, GA evolves towards a fewer number of request messages and shorter distance and larger number of tasks. Eventually, a string with the largest fitness value in population is selected. And after genetic\_operation is performed, the request messages are transferred to processors corresponding to bits set "1" in selected string.

In the receiver-initiated load sharing approach, a receiver finds a suitable sender to obtain a task. Therefore, a receiver performs GA-based load sharing algorithm. Processors received the request message from the receiver send accept message or reject message depending on its own CPU queue length.

/\* GA-based receiver-initiated load sharing algorithm \*/
Procedure Genetic\_algorithm Approach

```
{ Initialization()
    while (Check_load())
    if (Loadi <= T<sub>low</sub>) {
        Individual_evaluation();
        Genetic_operation();
        Message_evaluation();
    }
    Process a task in local processor;
}

Procedure Genetic_operation()

{ Local_improvement_operation();
    Reproduction();    Crossover(); }
```

#### 3. ADAPTIVE LOAD SHARING APPROACH

The purpose of this paper eventually uses the GA-based sender-initiated and receiver-initiated algorithm for new adaptive load sharing algorithm. The adaptive load sharing algorithm will be performed as the following. If the load of distributed system is low state, it is used conventional sender-initiated algorithm when a task insert to a specific processor. However, it is used a GA-based receiver-initiated load sharing algorithm when a task go out from a specific processor.

If the load of distributed system is heavy state, it is used a GA-based sender-initiated when a task insert to a specific processor. But it is use a conventional receiver-initiated load sharing algorithm when a task go out from a specific processor. The algorithm for above described contents is as following.

/\* An Adaptive Load sharing Algorithm \*/

```
If (load of distributed system < T_{low})
If (a specific processor == sender)
```

Use a conventional sender-initiated load sharing algorithm Else use a proposed a GA-based receiver-initiated load sharing algorithm

Else

If (a specific processor == receiver)

Use a conventional receiver-initiated load sharing algorithm Else use a proposed a GA-based sender-initiated load sharing algorithm

## 4. EXPERIMENTS

We executed several experiments on the proposed genetic algorithm approach to compare with the conventional sender-initiated and receiver-initiated algorithm. Our experiments have the following assumptions. First, each task size and task type are the same. Second, the number of parts(p) in a string is five. In genetic algorithm, crossover probability( $P_c$ ) is 0.7, mutation probability( $P_m$ ) is 0.05. The values of these

parameters  $P_c$ ,  $P_m$  were known as the most suitable values in various applications[3]. The table 2 shows the detailed contents of parameters used in our experiments.

Table 2. experimental parameters

number of processor	30
$P_c$	0.7
$P_m$	0.05
number of strings	50

## 4.1 Experiment of Sender-based Load Sharing Approach

Our experiments for sender-initiated load sharing approach have the following assumptions. The load rating over the systems is about 60 percent. The number of tasks to be performed is 3000. The weight for *TMP* is 0.025. The weight for *TMT* is 0.01. The weight for *TTP* is 0.02.

**[Experiment 1]** We compared the performance of proposed method with a conventional method in this experiment by using the parameters on the table 2. The experiment is to observe the change of response time when the number of tasks to be performed is 3000.

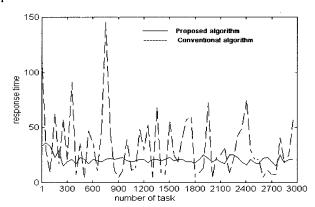


Fig 1. Result of response time

Fig. 1 shows result of the experiment 1. In conventional method, when the sender determines a suitable receiver, it select a processor in distributed systems randomly, and receive the load state information from the selected processor. The algorithm determines the selected processor as receiver if the load of randomly selected processor is  $T_{low}$  (light-load). These processes are repeated until a suitable receiver is searched. So, the result of response time shows the severe fluctuation. In the proposed algorithm, the algorithm shows the low response time activity load sharing the performs genetic operation considering the load state when it determines a receiver.

**[Experiment 2]** These experiments are to observe the performance when the probability of crossover is changed and is to observe the performance when the probability of mutation is changed.

Fig 2 shows the result of response time depending on the changes of  $P_c$  when  $P_m$  is 0.05. It shows high performance respectively when  $P_c$  is 0.7. Fig 3 shows the result of the

response time depending on the changes of  $P_m$  when  $P_c$  is 0.7. The result shows high performance respectively when  $P_m$  is 0.1.

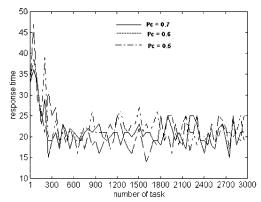


Fig 2. Result depending on the changes of  $P_c$ 

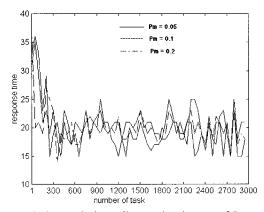


Fig 3. Result depending on the changes of  $P_m$ 

## 4.2 Experiment of Receiver-based Load Sharing Approach

Our experiments for receiver-initiated load sharing approach have the following assumptions. The load rating over the systems is about 20 percent. The number of tasks to be performed is 2000. The weight for *TMP* is 0.025. The weight for *TMT* is 0.01. The weight for *TTP* is 0.02.

[Experiment 3] We compared the performance of proposed method with a conventional method in this experiment by using the parameters on the table 2. The experiment is to observe the change of response time when the number of tasks to be performed is 2000.

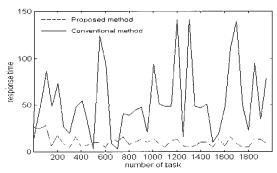


Fig 4. Result of response time

Through these several experiments, we verify that performance of the proposed scheme is better than that of the conventional scheme. The performance depending on the changes for parameters  $P_c$  and  $P_m$  is excellent than that of the conventional method.

#### 5. CONCLUSIONS

We propose a new adaptive load sharing scheme in distributed system that is based on the genetic algorithm. The genetic algorithm is used to decide to suitable candidate senders or receiver receivers which task transfer request messages should be sent off. Several experiments have been done to compare the proposed scheme with a conventional algorithm. Through the various experiments, performances of proposed scheme are better than those of conventional scheme on the response time and mean response time. However the proposed algorithm is sensitive to the weight values of *TMP*, *TMT* and *TTP*. As a further research, a study on method for releasing sensitivity of weight values is left.

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