

CLOSURE OPERATORS ON BL-ALGEBRAS

JUNG MI KO AND YONG CHAN KIM

ABSTRACT. We study relationships between closure operators and BL-algebras. We investigate the properties of closure operators and BL-homomorphisms on BL-algebras. We show that the image of a closure operator on a BL-algebra is isomorphic to a quotient BL-algebra.

1. Introduction and preliminaries

Hájek [10] introduced a BL-algebra as the foundation of the algebraic structures of fuzzy logic. Closure operators play an important role in (fuzzy) topological spaces [14, 16, 18], lattices [2-4, 6, 9, 11, 19, 20], Boolean algebras [2, 5, 6, 9, 11, 17] and deductive systems [1-3, 5, 6, 13, 20]. Recently, they have been developed in many view points [1, 3, 4-8, 10-20].

In this paper, we will study relationships between BL-algebras and closure operators in a sense Chakraborty [4]. We investigate the properties of closure operators and BL-homomorphisms on BL-algebras. Moreover, we study the first isomorphism theorem induced by closure operators on BL-algebras. We study the properties of quotient induced by closure operators on BL-algebras. We give the examples of them.

DEFINITION 1.1 ([5, 10, 20]). A lattice $(L, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ is called a *BL-algebra* if it satisfies the following conditions: for each $x, y, z \in L$,

- (B1) $(L, \odot, 1)$ is a commutative monoid,
- (B2) if $x \leq y$, then $x \odot z \leq y \odot z$ (\odot is an isotone operation),
- (B3) (Galois correspondence) $x \odot y \leq z$ if and only if $x \leq y \rightarrow z$,
- (B4) $x \wedge y = x \odot (x \rightarrow y)$,

Received July 30, 2003.

2000 Mathematics Subject Classification: 03B52, 94D05, 06D35.

Key words and phrases: closure operators, BL-algebras, deductive systems, BL-homomorphism, quotient BL-algebras.

$$(B5) \quad x \vee y = \left((x \rightarrow y) \rightarrow y \right) \wedge \left((y \rightarrow x) \rightarrow x \right),$$

$$(B6) \quad (x \rightarrow y) \vee (y \rightarrow x) = 1.$$

In a BL-algebra L , $x^* = (x \rightarrow 0)$ is called *complement* of $x \in L$.

LEMMA 1.2 ([10, 20]). In a BL-algebra $(L, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$, we have the following properties: for $x, y, z \in L$,

$$(1) \quad x = 1 \rightarrow x,$$

$$(2) \quad 1 = x \rightarrow x,$$

$$(3) \quad x \odot y \leq x, y,$$

$$(4) \quad x \odot y \leq x \wedge y,$$

$$(5) \quad y \leq x \rightarrow y,$$

$$(6) \quad x \odot y \leq x \rightarrow y,$$

$$(7) \quad x \leq y \text{ if and only if } 1 = x \rightarrow y,$$

$$(8) \quad x = y \text{ if and only if } 1 = x \rightarrow y = y \rightarrow x,$$

$$(9) \quad x \odot (x \rightarrow y) \leq y,$$

$$(10) \quad x \rightarrow (y \rightarrow z) = y \rightarrow (x \rightarrow z).$$

DEFINITION 1.3 ([5, 20]). Let L be a BL-algebra. A subset D of L is called a *deductive system* of L if it satisfies the following conditions:

$$(1) \quad 1 \in D,$$

$$(2) \quad \text{if } x, x \rightarrow y \in D, \text{ then } y \in D.$$

DEFINITION 1.4 ([5, 20]). Let \sim be an equivalence relation on A . Let $f : A^m \rightarrow A$ be an m -ary operation on A . We say that \sim is a *congruence* with respect to f if $a_i \sim b_i$ for each $i = 1, \dots, m$, then $f(a_1, \dots, a_m) \sim f(b_1, \dots, b_m)$.

THEOREM 1.5 ([5, 20]). If \sim is a congruence relation on a BL-algebra L . then $D = \{a \in L \mid a \sim 1\}$ is a deductive system.

THEOREM 1.6 ([5, 20]). Let L be a BL-algebra and D a deductive system of L . Define $a \sim b$ if and only if $(a \rightarrow b) \odot (b \rightarrow a) \in D$. Then \sim is a congruence relation with respect to $\rightarrow, \odot, *, \vee, \wedge$.

THEOREM 1.7 ([5, 20]). Let D be a deductive system of a BL-algebra L . Define on L/D which is the set of equivalence classes $\{|a| \mid a \in L\}$, for all $a, b \in L$,

$$|a| \leq |b| \text{ if and only if } a \rightarrow b \in D.$$

Then

$$(L/D, \leq, \wedge, \vee, \odot, \rightarrow, |0|, |1|)$$

is a BL-algebra where $|a| \wedge |b| = |a \wedge b|$, $|a| \vee |b| = |a \vee b|$, $|a| \odot |b| = |a \odot b|$, $|a| \rightarrow |b| = |a \rightarrow b|$.

DEFINITION 1.8 ([5, 20]). Let L, K be two BL-algebras. A map $h : L \rightarrow K$ is called a *BL-homomorphism* if for all $x, y \in L$, it satisfies the following conditions:

- (1) $h(x \rightarrow y) = h(x) \rightarrow h(y)$,
- (2) $h(x \odot y) = h(x) \odot h(y)$, $h(0) = 0$.

A BL-homomorphism $h : L \rightarrow K$ is called a *BL-isomorphism* if h is bijective.

THEOREM 1.9 ([5, 20]). Let L, K be two BL-algebras. Let $h : L \rightarrow K$ be a BL-homomorphism. Then, for all $x, y \in L$,

- (1) $h(x^*) = h(x)^*$, $h(1) = 1$,
- (2) if $x \leq y$, then $h(x) \leq h(y)$,
- (3) $h(x \wedge y) = h(x) \wedge h(y)$, $h(x \vee y) = h(x) \vee h(y)$,
- (4) if D is a deductive system of L , then $h(D)$ is a deductive system of K .

2. Closure operators and BL-algebras

DEFINITION 2.1 ([4]). Let $(L, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ be a BL-algebra. A map $c : L \rightarrow L$ is called a *closure operator* if for all $x, y \in L$, it satisfies the following conditions:

- (c1) $x \leq c(x)$.
- (c2) If $x \leq y$, then $c(x) \leq c(y)$.
- (c3) $c(c(x)) = c(x)$.
- (c4) $c(x) \odot c(y) \leq c(x \odot y)$.

THEOREM 2.2. Let $(L, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ be a BL-algebra. If $c : L \rightarrow L$ is a closure operator, then $D = \{a \in L \mid c(a) = 1\}$ is a deductive system.

PROOF. By (c1), $1 \in D$. Let $a, a \rightarrow b \in D$. Then $c(a) = 1, c(a \rightarrow b) = 1$. Since $b \geq a \odot (a \rightarrow b)$ from Lemma 1.2(9), by (c2) and (c4), $c(b) \geq c(a \odot (a \rightarrow b)) \geq c(a) \odot c(a \rightarrow b) = 1$. Thus $c(b) = 1$. So, $b \in D$. \square

THEOREM 2.3. Let $(L, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ be a BL-algebra, $c : L \rightarrow L$ a closure operator and $c(L) = \{x \in L \mid x = c(x)\}$. Then we have the following properties.

- (1) $c(L) = \{c(x) \mid x \in L\}$ and it is closed under \wedge and \rightarrow .

(2) $(c(L), \leq_c, \wedge_c, \vee_c, \odot_c, \rightarrow_c, c(0), 1)$ is a BL-algebra defined as follows: for each $x, y, z \in c(L)$,

$$\begin{aligned} x \wedge_c y &= c(x \wedge y), & x \vee_c y &= c(x \vee y), \\ x \odot_c y &= c(x \odot y), & x \rightarrow_c y &= c(x \rightarrow y). \end{aligned}$$

Furthermore,

$$\wedge_c = \wedge, \vee_c = \vee, \rightarrow_c = \rightarrow.$$

(3) If L satisfies $x \odot x = x$ for each $x \in L$, then

$$c(x \odot y) = c(x) \odot c(y).$$

PROOF. (1) We easily prove $c(L) = \{c(x) \mid x \in L\}$ from (c3). For each $x, y \in c(L)$, $x \wedge y \in c(L)$ from:

$$c(x) \wedge c(y) = x \wedge y \leq c(x \wedge y) \leq c(x) \wedge c(y) = x \wedge y.$$

We show that, for all $x, y \in c(L)$, $c(x \rightarrow y) = x \rightarrow y$. By (c1), $(x \rightarrow y) \leq c(x \rightarrow y)$. Since $c(x) \odot c(x \rightarrow y) \leq c(x \odot (x \rightarrow y)) \leq c(y)$ from (c4) and Lemma 1.2(9), then, by (B3), $c(x \rightarrow y) \leq c(x) \rightarrow c(y) = x \rightarrow y$.

(2) (B4-B5) We show that, for all $x, y, z \in c(L)$,

$$\begin{aligned} x \wedge_c y &= c(x \wedge y) = c(x \odot (x \rightarrow y)) = x \odot_c (x \rightarrow_c y) \\ x \vee_c y &= c(x \vee y) = c(((x \rightarrow y) \rightarrow y) \wedge ((y \rightarrow x) \rightarrow x)) \\ &= ((x \rightarrow_c y) \rightarrow_c y) \wedge_c ((y \rightarrow_c x) \rightarrow_c x). \end{aligned}$$

Furthermore, by (1),

$$\begin{aligned} x \wedge_c y &= c(x \wedge y) = x \wedge y, \\ x \vee_c y &= c(x \vee y) \\ &= c(((x \rightarrow y) \rightarrow y) \wedge ((y \rightarrow x) \rightarrow x)) \\ &= ((c(x) \rightarrow c(y)) \rightarrow c(y)) \wedge ((c(y) \rightarrow c(x)) \rightarrow c(x)) \\ &= c(x) \vee c(y) = x \vee y. \end{aligned}$$

Since $x \wedge_c y = x \wedge y$ and $x \vee_c y = x \vee y$ for all $x, y, z \in c(L)$, we have $\wedge_c = \wedge, \vee_c = \vee$ and $\leq_c = \leq$.

We substitute \leq for \leq_c in following statements.

(B3) We show that, for all $x, y, z \in c(L)$,

$$x \odot_c y \leq z \quad \text{if and only if} \quad x \leq y \rightarrow_c z.$$

(\Rightarrow) Since $c(x) \odot c(y) \leq c(x \odot y) = x \odot_c y \leq z$, we have $x = c(x) \leq c(y) \rightarrow c(z) = c(y \rightarrow z) = y \rightarrow_c z$.

(\Leftarrow) Since $x \odot_c y \leq (y \rightarrow_c z) \odot_c y = c(y \odot (y \rightarrow z)) \leq c(z)$, we have $x \odot_c y \leq z$.

(B1) $(c(L), \odot_c, 1)$ is a commutative monoid because, for all $x, y, z \in c(L)$,

$$\begin{aligned} x \odot_c y &= c(x \odot y) = c(y \odot x) = y \odot_c x \\ x \odot_c 1 &= c(x \odot 1) = c(x) = x. \end{aligned}$$

Suppose that there exist $x, y, z \in c(L)$ such that

$$x \odot_c (y \odot_c z) \not\leq (x \odot_c y) \odot_c z.$$

Then there exist $t \in c(L)$ such that

$$x \odot_c (y \odot_c z) \leq t < (x \odot_c y) \odot_c z.$$

By (B2),

$$\begin{aligned} x \odot_c (y \odot_c z) &\leq t \\ \text{if and only if} & \quad y \odot_c z \leq x \rightarrow_c t = c(x \rightarrow t) \\ \text{if and only if} & \quad y \leq z \rightarrow_c c(x \rightarrow t) \\ \text{if and only if} & \quad y \leq c(z \rightarrow c(x \rightarrow t)) \\ \text{if and only if} & \quad y \leq c(z) \rightarrow c(x \rightarrow t) \\ \text{if and only if} & \quad y \leq c(z) \rightarrow (c(x) \rightarrow c(t)) \\ \text{if and only if} & \quad y \leq c(x) \rightarrow (c(z) \rightarrow c(t)) \quad (\text{by Lemma 1.2(10)}) \\ \text{if and only if} & \quad y \leq x \rightarrow_c (z \rightarrow_c t) \\ \text{if and only if} & \quad x \odot_c y \leq z \rightarrow_c t \\ \text{if and only if} & \quad (x \odot_c y) \odot_c z \leq t. \end{aligned}$$

It is a contradiction. Hence $x \odot_c (y \odot_c z) \geq (x \odot_c y) \odot_c z$ for all $x, y, z \in c(L)$. Similarly, we have $x \odot_c (y \odot_c z) \leq (x \odot_c y) \odot_c z$ for all $x, y, z \in c(L)$.

(B2) If $x \leq y$ for each $x, y \in c(L)$,

$$x \odot_c z = c(x \odot z) \leq c(y \odot z) = y \odot_c z.$$

(B6) For each $x, y, z \in c(L)$,

$$\begin{aligned} (x \rightarrow_c y) \vee_c (y \rightarrow_c x) &= c\left(c(x \rightarrow y) \vee c(y \rightarrow x)\right) \\ &= c\left((c(x) \rightarrow c(y)) \vee (c(y) \rightarrow c(x))\right) \\ &= c(1) = 1. \end{aligned}$$

Thus $(c(L), \leq_c, \wedge_c, \vee_c, \odot_c, \rightarrow_c, c(0), 1)$ is a BL-algebra such that $\wedge_c = \wedge, \vee_c = \vee$ and $\rightarrow_c = \rightarrow$.

(3) Since $x \odot y \leq x \odot 1 = x$ and $x \odot y \leq y$, $c(x \odot y) \leq c(x)$ and $c(x \odot y) \leq c(y)$ imply

$$c(x \odot y) = c(x \odot y) \odot c(x \odot y) \leq c(x) \odot c(y).$$

□

EXAMPLE 2.4. Let $I = [0, 1]$ be the unit interval and $(I, \leq, \wedge, \vee, 0, 1)$ a lattice. We define $c : I \rightarrow I$ as follows:

$$c(x) = \begin{cases} \frac{1}{3}, & 0 \leq x \leq \frac{1}{3} \\ \frac{1}{2}, & \frac{1}{3} < x \leq \frac{1}{2} \\ 1, & \frac{1}{2} < x \leq 1. \end{cases}$$

(1) Define binary operations \odot and \rightarrow on I by

$$x \odot y = x \wedge y,$$

$$x \rightarrow y = \begin{cases} 1, & x \leq y \\ y, & \text{otherwise.} \end{cases}$$

Then $(I, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ is a BL-algebra (*Gödel logic*, ref. [10]). Since $c(x \odot y) = c(x) \odot c(y)$, c is a closure operator. From Theorem 2.3, $(c(I), \leq, \wedge, \vee, \odot, \rightarrow, \frac{1}{3}, 1)$ is a BL-algebra.

(2) Define binary operations \otimes and \Rightarrow on I by

$$x \otimes y = \max\{0, x + y - 1\},$$

$$x \Rightarrow y = \min\{1, 1 - x + y\}.$$

Then $(I, \leq, \wedge, \vee, \otimes, \Rightarrow, 0, 1)$ is a BL-algebra (*Lukasiewicz logic*, ref.[10]). But c is not a closure operator because

$$\frac{1}{3} = c\left(\frac{3}{7} \otimes \frac{4}{7}\right) \not\geq (c\left(\frac{3}{7}\right) \otimes c\left(\frac{4}{7}\right)) = \frac{1}{2} \otimes 1 = \frac{1}{2}.$$

Moreover, $c(I) = \{\frac{1}{3}, \frac{1}{2}, 1\}$ is not closed under \Rightarrow from

$$\frac{5}{6} = (c\left(\frac{1}{2}\right) \Rightarrow c\left(\frac{1}{3}\right)) \neq c\left(\frac{1}{2} \Rightarrow \frac{1}{3}\right) = c\left(\frac{5}{6}\right) = 1.$$

The converse of Theorem 2.3 cannot be true from the following example.

EXAMPLE 2.5. Let $I = [0, 1]$ be the unit interval and $(I, \leq, \wedge, \vee, 0, 1)$ a lattice. We define $c : I \rightarrow I$ as follows:

$$c(x) = \begin{cases} 0, & x = 0 \\ \frac{1}{4}, & 0 < x \leq \frac{1}{4} \\ \frac{1}{2}, & \frac{1}{4} < x \leq \frac{1}{2} \\ \frac{3}{4}, & \frac{1}{2} < x \leq \frac{3}{4} \\ 1, & \frac{3}{4} < x \leq 1. \end{cases}$$

Define binary operations \otimes and \Rightarrow on I by

$$x \otimes y = \max\{0, x + y - 1\},$$

$$x \Rightarrow y = \min\{1, 1 - x + y\}.$$

Then $(I, \leq, \wedge, \vee, \otimes, \Rightarrow, 0, 1)$ is a BL-algebra. We define $\wedge_c, \vee_c, \odot_c, \rightarrow_c$ as in Theorem 2.3(2). Also, $(c(L), \wedge_c = \wedge, \vee_c = \vee, \odot_c = \otimes, \rightarrow_c = \Rightarrow, 0, 1)$ is a BL-algebra where $c(L) = \{0, \frac{1}{4}, \frac{2}{4}, \frac{3}{4}, 1\}$ from the following statement.

If $\frac{k}{4} + \frac{l}{4} + \frac{m}{4} \leq 2$, then

$$\frac{k}{4} \otimes \left(\frac{l}{4} \otimes \frac{m}{4}\right) = \left(\frac{k}{4} \otimes \frac{l}{4}\right) \otimes \frac{m}{4} = 0.$$

If $\frac{k}{4} + \frac{l}{4} + \frac{m}{4} > 2$, then

$$\frac{k}{4} \otimes \left(\frac{l}{4} \otimes \frac{m}{4}\right) = \left(\frac{k}{4} \otimes \frac{l}{4}\right) \otimes \frac{m}{4} = \frac{k}{4} + \frac{l}{4} + \frac{m}{4} - 2.$$

Hence $x \otimes (y \otimes z) = (x \otimes y) \otimes z$ for all $x, y, z \in c(L)$. For all $x, y, z \in c(L)$.

$x \otimes y \leq z$ if and only if $x + y - 1 \leq z$ if and only if $x \leq 1 - y + z$
if and only if $x \leq \min\{1, 1 - y + z\}$ if and only if $x \leq y \Rightarrow z$.

If $x \leq y$, then

$$\begin{aligned} x &= x \wedge y = x \otimes 1 = x \otimes (x \rightarrow y), \\ ((x \rightarrow y) \rightarrow y) \wedge ((y \rightarrow x) \rightarrow x) &= (1 \rightarrow y) \wedge y = y = x \vee y, \\ (x \rightarrow y) \vee (y \rightarrow x) &= 1 \vee (1 - y + x) = 1. \end{aligned}$$

If $x > y$, it is similarly proved. Thus, it satisfies the conditions (B4-B6). Other conditions easily proved. But c is not a closure operator from

$$0 = c\left(\frac{3}{7} \otimes \frac{4}{7}\right) \not\geq c\left(\frac{3}{7}\right) \otimes c\left(\frac{4}{7}\right) = \frac{1}{2} \otimes \frac{3}{4} = \frac{1}{4}.$$

THEOREM 2.6. *Let $(L, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ and $(c(L), \leq_c, \wedge_c, \vee_c, \odot_c, \rightarrow_c, c(0), 1)$ be BL-algebras and let $c : L \rightarrow L$ be a closure operator. Then we have the following properties.*

- (1) $c : L \rightarrow c(L)$ is a BL-homomorphism.
- (2) If $D = \{a \in L \mid c(a) = 1\}$, then the map $\bar{c} : L/D \rightarrow c(L)$ defined by $\bar{c}(|a|) = c(a)$ is a BL-isomorphism.

PROOF. (1) It is easy from $c(x \odot y) = c(x) \odot_c c(y)$ and $c(x \rightarrow y) = c(x) \rightarrow_c c(y)$.

(2) From Theorem 2.2, D is a deductive system. Let $a \sim_D b$. Then $(a \rightarrow b) \odot (b \rightarrow a) \in D$. It implies

$$c\left((a \rightarrow b) \odot (b \rightarrow a)\right) = \left(c(a) \rightarrow c(b)\right) \odot_c \left(c(b) \rightarrow c(a)\right) = 1.$$

By Lemma 1.2(3,8), $c(a) = c(b)$. Thus, \bar{c} is well defined.

Since $c : L \rightarrow c(L)$ is a BL-homomorphism, $\bar{c} : L/D \rightarrow c(L)$ is a BL-homomorphism from the following statements:

$$\bar{c}(|x| \rightarrow |y|) = \bar{c}(|x \rightarrow y|) = c(x \rightarrow y) = c(x) \rightarrow_c c(y) = \bar{c}(|x|) \rightarrow_c \bar{c}(|y|),$$

$$\begin{aligned}\bar{c}(|x| \odot |y|) &= \bar{c}(|x \odot y|) = c(x \odot y) = c(x) \odot_c c(y) = \bar{c}(|x|) \odot_c \bar{c}(|y|), \\ \bar{c}(|0|) &= c(0).\end{aligned}$$

We only show that \bar{c} is bijective. Let $c(a) = c(b)$. From Lemma 1.2(8),

$$c(a) \rightarrow_c c(b) = c(b) \rightarrow_c c(a) = 1.$$

It implies

$$(c(a) \rightarrow_c c(b)) \odot_c (c(b) \rightarrow_c c(a)) = c((a \rightarrow b) \odot (b \rightarrow a)) = 1.$$

Then $(a \rightarrow b) \odot (b \rightarrow a) \in D$. Thus, $a \sim_D b$, that is, $|a| = |b|$. Hence \bar{c} is injective. Since c is surjective, \bar{c} is surjective. \square

We easily prove the following corollary.

COROLLARY 2.7. *Let X be a nonempty set and $P(X)$ a family of all subsets of X . Then $(P(X), \subset, \cap, \cup, \emptyset, X)$ is a lattice. For each $A, B \in P(X)$, we define the operations \odot and \rightarrow by*

$$A \odot B = A \cap B, \quad A \rightarrow B = A^c \cup B.$$

Then $(P(X), \subset, \cap, \cup, \odot, \rightarrow, \emptyset, X)$ is a BL-algebra.

EXAMPLE 2.8. Let $X = \{x_1, x_2, x_3\}$ be a set.

(1) Define $c : P(X) \rightarrow P(X)$ as follows:

$$c(A) = \begin{cases} \{x_1, x_2\}, & A \subset \{x_1, x_2\}, \\ X, & \text{otherwise.} \end{cases}$$

It satisfies the following conditions: for each $A, B \in P(X)$,

$$\begin{aligned}c(A \cap B) &= c(A) \cap c(B), \quad c(A \cup B) = c(A) \cup c(B), \\ c(A^c) &= c(A)^c \cup c(\emptyset).\end{aligned}$$

Since $A \rightarrow B = A^c \cup B$,

$$c(A \rightarrow B) = c(A^c \cup B) = c(A)^c \cup c(\emptyset) \cup c(B) = c(A) \rightarrow c(B).$$

From Theorem 2.3(2), $(c(P(X)), \subset, \cap, \cup, \odot = \cap, \rightarrow, c(\emptyset), X)$ is a BL-algebra. Hence $c : P(X) \rightarrow c(P(X))$ is a BL-homomorphism. Note that,

$$D = \{A \in P(X) \mid c(A) = X\} = \{\{x_3\}, \{x_1, x_3\}, \{x_2, x_3\}, X\}.$$

From Theorem 1.6, since

$$\begin{aligned} A \sim B & \text{ if and only if } (A \rightarrow B) \odot (B \rightarrow A) \in D \\ & \text{ if and only if } (A^c \cup B) \cap (B^c \cup A) \in D \\ & \text{ if and only if } (A \cap B) \cup (A \cup B)^c \in D. \end{aligned}$$

We can obtain:

$$\begin{aligned} \emptyset & \sim \{x_1\} \sim \{x_2\} \sim \{x_1, x_2\}, \\ X & \sim \{x_3\} \sim \{x_1, x_3\} \sim \{x_2, x_3\}. \end{aligned}$$

We obtain $P(X)/D = \{|\{x_1, x_2\}|, |X|\}$ and so $\bar{c} : P(X)/D \rightarrow c(P(X))$, $|A| \mapsto c(A)$, is a BL-isomorphism by Theorem 2.6(2).

(2) Define $c_1 : P(X) \rightarrow P(X)$ as follows:

$$c_1(A) = \begin{cases} \{x_1\}, & A \subset \{x_1\}, \\ X, & \text{otherwise.} \end{cases}$$

Then c_1 is not a closure operator because

$$X = c_1(\{x_1, x_2\}) \cap c_1(\{x_1, x_3\}) \not\subseteq c_1(\{x_1, x_2\} \cap \{x_1, x_3\}) = \{x_1\}.$$

(3) We define $(I, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ and a closure operator c as same as in Example 2.4(1). Then $c : I \rightarrow c(I)$ is a BL-homomorphism and $D = \{\frac{1}{2} < x \leq 1 \mid c(x) = 1\}$. From Theorem 1.6, we can obtain:

$$\begin{aligned} \frac{1}{3} & \sim x, & \forall x \in [0, \frac{1}{3}], \\ \frac{1}{2} & \sim y, & \forall y \in (\frac{1}{3}, \frac{1}{2}], \\ 1 & \sim z, & \forall z \in (\frac{1}{2}, 1]. \end{aligned}$$

We obtain $I/D = \{|\frac{1}{3}|, |\frac{1}{2}|, |1|\}$. Define $\bar{c} : I/D \rightarrow c(I)$ by $\bar{c}(|x|) = c(x)$. By Theorem 2.6, \bar{c} is a BL-isomorphism.

THEOREM 2.9. *Let L, K be two BL-algebras and $h : L \rightarrow K$ a BL-homomorphism. Let $c_1 : L \rightarrow c(L)$ and $c_2 : K \rightarrow c(K)$ be closure operators. Let $D_1 = \{x \in L \mid c_1(x) = 1\}$ and $D_2 = \{y \in K \mid c_2(y) = 1\}$. Then we have the following properties.*

(1) *If $h(c_1(x)) \leq c_2(h(x))$ for all $x \in L$, which is called a closed map, a map $\bar{h} : L/D_1 \rightarrow K/D_2$ defined by $\bar{h}(|x|) = |h(x)|$ is a BL-homomorphism.*

(2) *If h is surjective such that $h(c_1(x)) = c_2(h(x))$ for all $x \in L$ and $h(c_1(x)) = 1$ implies $c_1(x) = 1$, then a map $\bar{h} : L/D_1 \rightarrow K/D_2$ is a BL-isomorphism.*

PROOF. (1) Let $a \sim_{D_1} b$. From Theorem 1.6, $(a \rightarrow b) \odot (b \rightarrow a) \in D_1$. Since $c\left((a \rightarrow b) \odot (b \rightarrow a)\right) = 1$, by Theorem 1.9(1), we have $h\left(c\left((a \rightarrow b) \odot (b \rightarrow a)\right)\right) = 1$. Since $h(c_1(x)) \leq c_2(h(x))$, $c_2\left(h\left((a \rightarrow b) \odot (b \rightarrow a)\right)\right) = 1$. Then $h(a) \sim_{D_2} h(b)$. So, the map \bar{h} is well defined. Since $|x \rightarrow y| = |x| \rightarrow |y|$ and $|h(x) \rightarrow h(y)| = |h(x)| \rightarrow |h(y)|$ from Theorem 1.7, we have

$$\begin{aligned} \bar{h}(|x| \rightarrow |y|) &= \bar{h}(|x \rightarrow y|) = |h(x) \rightarrow h(y)| \\ &= |h(x)| \rightarrow |h(y)| = \bar{h}(|x|) \rightarrow \bar{h}(|y|). \end{aligned}$$

Similarly, $\bar{h}(|x| \odot |y|) = \bar{h}(|x|) \odot \bar{h}(|y|)$, $\bar{h}(|0|) = |0|$. Thus, \bar{h} is BL-homomorphism.

(2) Since h is surjective, \bar{h} is surjective. We only show that \bar{h} is injective. Let $h(x) \sim_{D_2} h(y)$. Then $(h(x) \rightarrow h(y)) \odot (h(y) \rightarrow h(x)) \in D_2$, that is,

$$\begin{aligned} 1 &= c_2\left((h(x) \rightarrow h(y)) \odot (h(y) \rightarrow h(x))\right) \\ &= \left(c_2(h(x)) \rightarrow_{c_2} c_2(h(y))\right) \odot_{c_2} \left(c_2(h(y)) \rightarrow_{c_2} c_2(h(x))\right). \end{aligned}$$

By Lemma 1.2 (3,8), $c_2(h(x)) = c_2(h(y))$. Since $c_2(h(x)) = h(c_1(x))$ for all $x \in X$, $h(c_1(x)) = h(c_1(y))$. It implies

$$\begin{aligned} &h\left(c_1\left((x \rightarrow y) \odot (y \rightarrow x)\right)\right) \\ &= \left(h(c_1(x)) \rightarrow h(c_1(y))\right) \odot \left(h(c_1(y)) \rightarrow h(c_1(x))\right) \\ &= 1 \odot 1 = 1. \end{aligned}$$

Hence $c_1\left((x \rightarrow y) \odot (y \rightarrow x)\right) = 1$. Thus, $(x \rightarrow y) \odot (y \rightarrow x) \in D_1$. Hence $x \sim_{D_1} y$. □

EXAMPLE 2.10. Let $X = \{x_1, x_2, x_3\}$ and $Y = \{y_1, y_2\}$ be two sets. Define $h : P(X) \rightarrow P(Y)$ as follows:

$$\begin{aligned} h(\emptyset) &= \emptyset, \quad h(X) = Y, \\ h(\{x_1\}) &= \{y_1\}, \quad h(\{x_2\}) = \{y_2\}, \quad h(\{x_3\}) = \emptyset, \\ h(\{x_1, x_2\}) &= \{y_1, y_2\}, \quad h(\{x_1, x_3\}) = \{y_1\}, \quad h(\{x_2, x_3\}) = \{y_2\}. \end{aligned}$$

It satisfies the following conditions: for each $A, B \in P(X)$,

$$h(A \cap B) = h(A) \cap h(B), \quad h(A \cup B) = h(A) \cup h(B), \quad h(A^c) = h(A)^c.$$

Since $A \rightarrow B = A^c \cup B$,

$$h(A \rightarrow B) = h(A^c \cup B) = h(A)^c \cup h(B) = h(A) \rightarrow h(B).$$

Hence $h : P(X) \rightarrow P(Y)$ is a BL-homomorphism.

Define $c_1 : P(X) \rightarrow P(X)$ as follows:

$$c_1(A) = \begin{cases} \{x_1, x_3\}, & A \subset \{x_1, x_3\}, \\ X, & \text{otherwise.} \end{cases}$$

Define $c_2 : P(Y) \rightarrow P(Y)$ as follows:

$$c_2(B) = \begin{cases} \{y_1\}, & B \subset \{y_1\}, \\ Y, & \text{otherwise.} \end{cases}$$

Then c_1 and c_2 are closure operators. It implies

$$h(c_1(A)) = \begin{cases} h(\{x_1, x_3\}) = \{y_1\} = c_2(h(A)), & A \subset \{x_1, x_3\}, \\ h(X) = Y = c_2(h(A)), & \text{otherwise.} \end{cases}$$

Moreover, $D_1 = \{\{x_2\}\{x_1, x_2\}\{x_2, x_3\}, X\}$, $D_2 = \{\{y_2\}, Y\}$. We define $\bar{h} : P(X)/D_1 \rightarrow P(Y)/D_2$ as $\bar{h}(|\{x_1, x_3\}|) = |\{y_1\}|$ and $\bar{h}(|X|) = |Y|$. Since $h(c_1(A)) = c_2(h(A))$ for each $A \subset X$ and $h(c_1(A)) = Y$ implies $c_1(A) = X$ by Theorem 2.9(2), \bar{h} is a BL-isomorphism.

EXAMPLE 2.11. Let $I = [0, 1]$ be the unit interval and $(I, \leq, \wedge, \vee, 0, 1)$ be a lattice. We define $c_1, c_2 : I \rightarrow I$ as follows:

$$c_1(x) = \begin{cases} \frac{1}{3}, & 0 \leq x \leq \frac{1}{3} \\ \frac{1}{2}, & \frac{1}{3} < x \leq \frac{1}{2} \\ 1, & \frac{1}{2} < x \leq 1, \end{cases} \quad c_2(x) = \begin{cases} \frac{1}{2}, & 0 \leq x \leq \frac{1}{2} \\ \frac{3}{4}, & \frac{1}{2} < x \leq \frac{3}{4} \\ 1, & \frac{3}{4} < x \leq 1. \end{cases}$$

Define on I binary operations \odot and \rightarrow by

$$x \odot y = x \wedge y,$$

$$x \rightarrow y = \begin{cases} 1, & x \leq y \\ y, & \text{otherwise.} \end{cases}$$

Then $(I, \leq, \wedge, \vee, \odot, \rightarrow, 0, 1)$ is a BL-algebra. Furthermore, for each $i = 1, 2$, c_i is a closure operator satisfying $c_i(x \odot y) = c_i(x) \odot c_i(y)$. By Theorem 2.3, $(c_1(I), \leq, \wedge, \vee, \odot, \rightarrow, \frac{1}{3}, 1)$ and $(c_2(I), \leq, \wedge, \vee, \odot, \rightarrow, \frac{1}{2}, 1)$ are BL-algebras. Define a function $h : I \rightarrow I$ by

$$h(x) = \begin{cases} \frac{3}{2}x, & 0 \leq x \leq \frac{1}{2} \\ \frac{1}{2}x + \frac{1}{2}, & \frac{1}{2} < x \leq 1. \end{cases}$$

We obtain $\bar{h} : I/D_1 \rightarrow I/D_2$ as

$$\bar{h}(|\frac{1}{3}|) = |\frac{1}{2}|, \quad \bar{h}(|\frac{1}{2}|) = |\frac{3}{4}|, \quad \bar{h}(|1|) = |1|.$$

Since $h(c_1(x)) = c_2(h(x))$ for each $x \in I$ and $h(c_1(x)) = 1$ implies $c_1(x) = 1$ by Theorem 2.9, \bar{h} is a BL-isomorphism.

References

- [1] L. Biacino and G. Gerla, *An extension principle for closure operators*, J. Math. Anal. Appl. **198** (1996), 1–24.
- [2] G. Birkhoff, *Lattice theory, 3rd Edition*, A.M.S, Rhode Island (1967).
- [3] N. Caspard and B. Monjardet, *The lattices of closure systems, closure operators, and implicational systems on a finite set: a survey*, Discrete Appl. Math. **127** (2003), 241–269.
- [4] M. K. Chakraborty and J. Sen, *MV-algebras embedded in a CL-algebra*, Internat. J. Approx. Reason. **18** (1998), 217–229.
- [5] A. DiNola, G. Georgescu and L. Leustean, *Boolean products of BL-algebras*, J. Math. Anal. Appl. **251** (2000), 106–131.
- [6] P. Flondor and M. Sularia, *On a class of residuated semilattice monoid*, Fuzzy Sets and Systems **138** (2003), 149–176.
- [7] G. Gerla, *An extension principle for fuzzy logics*, Math. Log. Quart. **40** (1994), 357–380.
- [8] G. Gerla and L. Scarpati, *Extension principles for fuzzy set theory*, Inform. Sci. **106** (1998), 49–69.
- [9] G. Grätzer, *Universal Algebra*, Springer-Verlag, New York (1978).
- [10] P. Hájek, *Metamathematics of Fuzzy Logic*, Kluwer Academic Publishers, Dordrecht (1998).
- [11] A. Higuchi, *Lattices of closure operators*, Discrete Math. **179** (1998), 267–272.
- [12] U. Höhle, *On the fundamentals of fuzzy set theory*, J. Math. Anal. Appl. **201** (1996), 786–826.
- [13] U. Höhle and S. E. Rodabaugh, *Mathematics of fuzzy sets*, Kluwer Academic Publishers (1999).
- [14] Y. C. Kim, *Initial L-fuzzy closure spaces*, Fuzzy Sets and Systems **133** (2003), 277–297.

- [15] J. M. Ko and Y. C. Kim, *Some properties of BL-algebras*, J. of Fuzzy Logic and Intelligent Systems **11(3)** (2001), 286–291.
- [16] A. S. Mashhour and M.H. Ghanim, *Fuzzy closure spaces*, J. Math. Anal. Appl. **106** (1985), 154–170.
- [17] F. Ranzato, *Pseudocomplements of closure operators on posets*, Discrete Math. **248** (2002), 143–155.
- [18] R. Srivastava and A. K. Srivastava and A. Choubey, *Fuzzy closure spaces*, J. Fuzzy Math. **2** (1994), 525–534.
- [19] E. Turunen, *Algebraic structures in fuzzy logic*, Fuzzy Sets and Systems **52** (1992), 181–188.
- [20] ———, *Mathematics behind fuzzy logic*, A Springer-Verlag Co., 1999.

Department of Mathematics
Kangnung National University
Gangwondo 210-702, Korea
E-mail: jmko@kangnung.ac.kr
yck@kangnung.ac.kr