# On The Category Of G-sets

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#### 1. Introduction.

Let G be a group and S a set. An operation of G on S is a map  $f: G \times S \longrightarrow S$  which, if we denote f(g,s) by gs for all g in G and s in S, satisfies;

- (i) 1s=s for all s in S and the identity element  $1 \in G$ ,
- (ii)  $(g_1g_2)s=g_1(g_2s)$  for all  $g_1,g_2 \in G$  ((1),(2)).

A set S together with an operation of G on S is called a G-set ((1)).

The category  $\mathcal{C}(G)$  whose object is the single group G and set of morphisms  $\operatorname{Hom}_{\mathcal{C}(G)}(G,G)$  is the set G is called the category of the group G.

Definition. Let G be a group. If  $S_1$  and  $S_2$  are G-sets, then a G-morphism from  $S_1$  to  $S_2$  is a map of the sets  $f: S_1 \rightarrow S_2$  satisfying

f(gs)=gf(s) for all g in G and s in S.

The category with objects the G-sets and morphisms the G-morphisms is called the *category* of G-sets and denoted by G-sets.

In this paper we shall prove the theorem: The category  $Funct(\mathcal{G}(G), Sects)$  of functors from the category of the group G to the category Sets of all sets is equivalent to the category G-sets.

#### 2. G-sets.

**Proposition.** (a) For each G-set S, id,  $: S \longrightarrow S$  is a G-morphism.

- (b) If  $S_1, S_2, S_3$  are G-sets and  $f: S_1 \longrightarrow S_2$  and  $h: S_2 \longrightarrow S_3$  are G-morphisms, then the ordinary composition of maps  $hf: S_1 \longrightarrow S_3$  is a G-morphism.
- (c) The following data define a category which is called the category of G-sels and is denoted by G-Sets.
  - (i) The objects of G-Sets are the G-sets.
  - (ii) For each pair of objects  $S_1$  and  $S_2$  of G-Sets, G-Sets( $S_1$ ,  $S_2$ ) is the set of all G-morphisms from  $S_1$  to  $S_2$ .
- (iii) For each triple  $S_1$ ,  $S_2$  and  $S_3$  of objects of G-Sets, the composition map  $(G_1, S_2) \times G_2 \times G_3 \times G_3 \times G_4 \times G_4 \times G_5 \times$

Proof. (a)  $id_s(gs)=gs=gid_s(s)$  for all g in G and s in S.

- (b) hf(gs)=h(gf(s))=g(h(f(s)))=g(hf(s)) for all g in G and s in S<sub>1</sub>.
- (c) If  $S_1, S_2, S_3, S_4$  are objects in G-sets and f is in G-sets  $(S_1, S_2)$ , g is in G-sects  $(S_2, S_3)$ , and h is in G-sets  $(S_3, S_4)$ , then h(gf) = (hg)f since the ordinary composition of maps is associative and h(gf) and (hg)f are G-morphisms by (a).

Next for each object S in G-sets, there is an id, in G-Sets (S, S) such that for each object S<sub>1</sub> in G-sets, we have  $fid_i = f$  for all f in G-sets (S, S<sub>1</sub>) while  $id_i g = g$  for all g in G-sets (S<sub>1</sub>, S), which completes the proof.

### 3. The category G-sets.

By the proposition in 2, we obtain the category G-sets with objects G-sets and arrow G-morphisms.

Lemma 1. Let G be a group and  $\theta(G)$  be the category of G. Let  $F: \theta(G) \longrightarrow Sets$  be a functor. If  $\alpha(F)$  is the set F(G) and for each g in G and s in F(G), f(g,s)=gs = F(g)(s), then f is an operation of G on F(G) and  $\alpha(F)$  is the G-set.

**Proof.** For the identity element  $1 \in G$  and for  $g_1$ ,  $g_2 \in G$ ,  $1s = F(1)(s) = 1_{F(G)}(s) = s$  for all s in F(G).

$$(g_1, g_2)s = F(g_1g_2)(s) = F(g_1)F(g_2)(s) = F(g_1) (F(g_2)(s)) = g_1(g_2s).$$

Lemma 2. Let  $F_1, F_2 : \mathcal{C}(G) \longrightarrow Sets$  be functors and let  $\rho : F_1 \longrightarrow F_2$  be a morphism of functors. Then the map  $\rho : \alpha(F_1) \longrightarrow \alpha(F_2)$   $\downarrow \qquad \qquad \downarrow \qquad$ 

is a morphism of G-sets.

**Proof.** For all g in G and s in  $\alpha(F_1)$ ,  $\rho_c(gs) = \rho_c(F_1(g)(s)) = F_2(g)(\rho_c(s)) = g(\rho_c(s))$ , since  $\rho$  is morphism of functors.

**Theorem.** Let Funct(G(G), Sets) be the category of functors from the category of the group G to the category Sets of all sets. For every functor  $F: G(G) \longrightarrow Sets$  and for all morphisms of functors  $\rho: F_1 \longrightarrow F_2$  in Funct G(G), Sets), if  $\alpha(F) = F(G)$  and  $\alpha(\rho) = \rho: F_1(G) \longrightarrow F_2(G)$ , then  $\alpha$  is an equivalence functor from the category Funct(G(G), Set(G(G)) to the category G-sets.

**Proof.** For two morphisms  $\rho_1: F_1 \longrightarrow F_2$  and  $\rho_2: F_2 \longrightarrow F_3$  in Funct ( $\mathcal{C}(G)$ , Sets),

$$\alpha(\rho_2\rho_1) = (\rho_2\rho_1)_G = (\rho_2)_G(\rho_1)_G = \alpha(\rho_2)\alpha(\rho_1)$$

and for the identity morphism  $1_F: F \longrightarrow F$ ,

$$\alpha(1_F) = (1_F)_G = 1_{F(G)} = 1_{\alpha(F)}$$

Therefore  $\alpha$  is a functor.

Next, for every G-sets S in Ob(G-sets) and all G-morphisms  $f: S_1 \longrightarrow S_2$  in G-sets, if  $\beta(S)(G)=S \in Ob(Sets)$  and  $\beta(f)_G=f$ , then for  $S_1 \xrightarrow{f_1} S_2 \xrightarrow{f_2} S_3$  in G-sets

$$(\beta(f_2f_1))_G = f_2f_1 = \beta(f_2)_G\beta(f_1)_G = (\beta(f_2)\beta(f_1))_G$$

and for  $1_s: S \longrightarrow S$ ,  $\beta(1_s)_c = 1_s = 1_{\beta(s) \in G}$ 

Hence  $\beta$  is also a functor from G-sets to Funct ( $\theta$ (G), Sets).

$$(\beta\alpha(F)) = \beta(\alpha(F)) = \beta(F(G)) \in Funct (\beta(G), Sets).$$

But  $(\beta\alpha(F))(G) = \beta(FG))(G) = F(G) = [{}^{1}_{Fund}(\sigma_{(G)}, Sets)}(F)](G)$ 

Therefore  $\beta \alpha = 1_{\text{Funct}(\delta(G), \text{ Sets})}$ 

Similarly  $\alpha\beta = 1_{S-sets}$ , which completes the proof.

## REFERENCES

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