

## *Inf-Preserving Functors from $\underline{A}$ to $Ens$*

CHANG KU IM

**1. Introduction.** Let  $I$  and  $\underline{A}$  be an index category and a small category respectively. For each object  $A$  of  $\underline{A}$  the constant diagram  $A_I: I \rightarrow \underline{A}$  is defined by  $A_I(i) = A$ ,  $A_I(k) = 1_A$  for each object  $i$  and map  $k: i \rightarrow j$  of  $I$ . A lower bound  $(A, u)$  of a diagram  $\Gamma: I \rightarrow \underline{A}$  consists of an object  $A$  of  $\underline{A}$  and a natural transformation  $u: A_I \rightarrow \Gamma$ . The lower bound  $(A, u)$  of  $\Gamma$  will be called the infimum of  $\Gamma$  if for every lower bound  $(A', u')$  of  $\Gamma$  there exists a unique map  $a: A' \rightarrow A$  such that  $u(i)a = u'(i)$  for all objects  $i$  of  $I$ , and we write it by  $\inf \Gamma = (A, u)$ . A functor  $F: \underline{A} \rightarrow \underline{B}$  is called *the inf-preserving functor* or we say that it *preserves the infimums* if for every diagram  $\Gamma: I \rightarrow \underline{A}$ ,  $\inf \Gamma = (A, u) \Rightarrow \inf (F\Gamma) = (F(A), F \cdot u)$ . An upper bound, the supremum of a diagram and the sup-preserving functors are also defined dually. We write the opposite category of  $\underline{A}$  and the category of sets by  $\underline{A}^\circ$  and  $Ens$  respectively. J. Lambek [1] proved that  $\underline{A}$  is embedded as a sup-dense subcategory into a sup-complete category  $\underline{A}'$  of all functors from  $\underline{A}^\circ$  to  $Ens$  and the embedding functor of  $\underline{A}$  into the category  $\underline{A}''$  of all inf-preserving functors from  $\underline{A}^\circ$  to  $Ens$  is sup-dense and sup-preserving. Further he proved that the category  $\underline{A}''$  is inf-complete. The purpose of this note is to prove that the opposite  $\underline{A}'''$  of the category  $\underline{A}''$  of all inf-preserving functors from  $\underline{A}$  to  $Ens$  is sup-complete and it is inf-complete if for any diagram  $\theta$  with  $\inf \theta = (z, t)$ ,  $t$  is a natural equivalence.

Throughout this note we assume that every diagram has the small index category.

**2. Inf-preserving functors.** Let  $\{o\}$  be a typical one element set and  $T: \underline{A} \rightarrow Ens$ . We may associate with the element  $x$  of  $T(A)$  for all  $A$  of  $\underline{A}$  the map  $\hat{x}: \{o\} \rightarrow T(A)$  such that  $\hat{x}(o) = x$ . The following lemmas will be stated whose proofs are to be found in [1] and [3] respectively.

**LEMMA 1.** *For any object  $A$  of  $\underline{A}$  the functors  $[A, \ ]: \underline{A} \rightarrow Ens$  and  $[ \ , A]: \underline{A}^\circ \rightarrow Ens$  preserves infimums.*

**LEMMA 2.** *Let  $T$  and  $T'$  be two functors from  $\underline{A}$  to  $\underline{B}$  and  $\eta: T \rightarrow T'$  a natural equivalence. Then  $T$  preserves infimums if and only if  $T'$  does.*

**PROPOSITION 1.** *The functor  $T: \underline{A} \rightarrow \underline{B}$  preserves infimums if and only if  $[B, T(\ )]: \underline{A} \rightarrow Ens$  preserves infimums for all  $B$  in  $\underline{B}$ .*

*Proof.* Assume that  $T$  preserves infimums. The functor  $[B, T(\ )]$  arises by composition from the inf-preserving functor  $T: \underline{A} \rightarrow \underline{B}$  and the functor  $[B, \ ]: \underline{B} \rightarrow Ens$ . But the functor  $[B, \ ]$  preserves infimum by the lemma 1. Hence it also preserves infimums. Conversely, assume that the functor  $[B, T(\ )]$  preserves infs for all  $B$  in  $\underline{B}$ . Let  $D: I \rightarrow \underline{A}$  be a diagram in  $\underline{A}$  with  $\inf D = (A, u)$ . Then the infimum of the diagram  $[B, T(\ )] \cdot D: I \rightarrow Ens$  is  $([B, T(A)], v)$ , where  $v(i) = [B, T(u(i))]$  for each  $i$  of  $I$ . Let  $t(i): B \rightarrow T(D(i))$  be natural in  $i$  of  $I$ . We associate with  $t(i)$  mapping  $\hat{t}(i): \{o\} \rightarrow [B, T(D(i))]$ . Hence there exists a unique map  $g: \{o\} \rightarrow [B, T(A)]$  such that  $t(i) = v(i) \cdot g$ . Since  $\hat{t}(i)(o) = v$

(i)  $\cdot (g(o))$ , we have a unique element  $g(o)$  of  $[B, T(A)]$  such that  $t(i) = v(i)g(o) = T(u(i)) \cdot g(o)$ . Hence  $\inf TD = [T(A), T \cdot u]$ .

PROPOSITION 2. Let  $T: \underline{A} \rightarrow \underline{C}$  be an embedding functor and  $\underline{C}_0$  be the subcategory of  $\underline{C}$ , consisting of all objects  $C$  in  $\underline{C}$  such that the functor  $[C, T(\ )]: \underline{A} \rightarrow \text{Ens}$  preserves infimums. Then

- (i) the image  $T(\underline{A})$  of  $T$  is contained in  $\underline{C}_0$ .
- (ii)  $\underline{C}_0$  is the largest subcategory of  $\underline{C}$  such that the induced embedding functor  $\underline{A} \rightarrow T(\underline{A}) \rightarrow \underline{C}_0$  preserves infimums.
- (iii) Every supremum of any diagram  $\Delta: J \rightarrow \underline{C}$  in  $\underline{C}$  is contained in  $\underline{C}_0$ .
- (iv) For any diagram  $\theta: K \rightarrow \underline{C}_0$  with  $\inf \theta = (C', t)$  in  $\underline{C}$  if  $t$  is a natural equivalence then  $C'$  is contained in  $\underline{C}_0$ .

*Proof.* (i) Since  $T$  is an embedding functor, we have a natural equivalence  $\mu: [A, \ ] \cong [T(A), T(\ )]$  for each object  $A$  of  $\underline{A}$ , where  $[A, \ ]$  and  $[T(A), T(\ )]$  are two functors from  $\underline{A}$  to  $\text{Ens}$ . The functor  $[A, \ ]$  preserves infimums by the lemma 1. Hence the functor  $[T(A), T(\ )]$  preserves infimums, by the lemma 2. Therefore  $T(\underline{A}) \in \underline{C}_0$  and  $T(\underline{A}) \subset \underline{C}_0$ .

(ii) Let  $\underline{A} \rightarrow T(\underline{A}) \rightarrow \underline{C}'$  be an inf-preserving induced embedding functor. By the proposition 1, for each object  $C'$  of  $\underline{C}'$ : the functor  $[C', T(\ )]$  preserves infimums. Hence  $C'$  is contained in  $\underline{C}_0$ .

(iii) Let  $\Delta: J \rightarrow \underline{C}_0$  be any diagram with  $\sup \Delta = (C, v)$  in  $\underline{C}$ . Assume that  $D: I \rightarrow \underline{A}$  is a diagram with  $\inf D = (A, u)$  then  $\inf T \cdot D = ((TA), T \cdot u)$  in  $\underline{C}_0$ . Let  $g: C_I \rightarrow TD$  be a natural transformation, then we have the map  $g(i)v(j)$  in  $\underline{C}$  such that

$$\begin{array}{ccc} \Delta(j) & \xrightarrow{v(j)} & C \\ & \searrow & \downarrow g(i) \\ & & TD(i) \end{array} \quad \text{commutes.}$$

Hence there exists a unique map  $S(j): \Delta(j) \rightarrow T(A)$  such that

$$\begin{array}{ccc} \Delta(j) & \xrightarrow{v(j)} & C \\ S(j) \downarrow & \searrow f & \downarrow g(i) \\ T(A) & \xrightarrow{T(u(i))} & TD(i) \end{array} \quad \text{commutes.}$$

Since  $T(u(i)) \cdot S(j): \Delta(j) \rightarrow TD(i)$  is a natural in  $i \in I$ , so does  $S(j)$  in  $j \in J$ . Hence there exists a unique map  $f: C \rightarrow T(A)$  such that  $S(j) = f \cdot v(j)$ . Therefore  $T(u(i)) \cdot S(j) = T(u(i)) \cdot f \cdot v(j) = g(i)v(j)$ , hence  $g(i) = T(u(i))f$ , that is,  $T(A)$  is the infimum in  $\underline{C}_0 \cup \{C\}$ , so that  $\underline{C}_0 \cup \{C\} = \underline{C}_0$  by (ii).

(iv) Consider any diagrams  $\theta: K \rightarrow \underline{C}_0$  with  $\inf \theta = (C', t)$  in  $\underline{C}$ , where  $t$  is a natural equivalence and  $D: I \rightarrow \underline{A}$  with  $\inf D = (A, u)$ . Then  $T(\underline{A}) \rightarrow TD(i)$  is a natural in  $i \in I$ . Let  $x(i): C' \rightarrow TD(i)$  be a natural in  $i \in I$ , for each  $k$  of  $K$  then we have a natural map  $x(i)t^{-1}(k): \theta(k) \rightarrow TD(i)$  in  $i$  of  $I$ . There exists a unique map  $y(k): \theta(k) \rightarrow T(A)$

such that  $T(u(i)) \cdot y(k) = x(i)t^{-1}(k)$ . Hence  $T(u(i)) \cdot y(k)t(k) = x(i)$ .

$$\begin{array}{ccc}
 \theta(k) & \xrightarrow{t^{-1}(k)} & C' \\
 \downarrow y(k) & & \downarrow x(i) \\
 T(A) & \xrightarrow{T(u(i))} & TD(i)
 \end{array} \quad \text{commutes.}$$

By (ii),  $C'$  must be in  $\mathcal{C}_0$ .

**3. Category of inf-preserving functors.** Let  $\underline{A}$  be any small category. We shall write  $[\underline{A}, \text{Ens}]_{\text{inf}}$  for the category of all inf-preserving functors from  $\underline{A}$  to  $\text{Ens}$ .

**LEMMA 3.** For each object  $C$  of  $\mathcal{C}$  and each functor  $S: \mathcal{C} \rightarrow \text{Ens}$ , there is a bijection  $\psi: S(c) \cong [h_c S]$ , where  $h_c$  is a functor  $[C, \ ]: \mathcal{C} \rightarrow \text{Ens}$  and the map  $\psi$  is defined for  $x \in S(c)$  and  $f \in h_c(A)$  for  $A$  of  $\mathcal{C}$  as  $(\psi(x))(f) = (S(f))(x)$ , [2].

By the lemma 1 the canonical embedding  $H: \underline{A} \rightarrow [\underline{A}, \text{Ens}]^\circ$  induce the embedding of  $\underline{A}$  into  $[\underline{A}, \text{Ens}]_{\text{inf}}$ . Using the proposition 2, we shall show the following theorem.

**THEOREM.** (1) The category  $[\underline{A}, \text{Ens}]_{\text{inf}}^\circ$  is sup-complete.

(2) For any diagram  $\theta$  with  $\text{inf } \theta = (x, t)$  in  $[\underline{A}, \text{Ens}]^\circ$  if  $t$  is a natural equivalence then  $[\underline{A}, \text{Ens}]_{\text{inf}}$  is inf-complete.

*Proof.* (1) Let  $\underline{B}$  be the category of all functors  $T: \underline{A} \rightarrow \text{Ens}$  in  $[\underline{A}, \text{Ens}]^\circ$  such that  $[T, H(\ )]$  preserves infimums. By the lemma 3,  $T \cong [T, H(\ )]$ . Therefore we have  $[\underline{A}, \text{Ens}]_{\text{inf}}^\circ = \underline{B}$  by the lemma 2. Hence the category  $\underline{B}$  is sup-complete by (iii) of the proposition 2.

(2) Since  $[\underline{A}, \text{Ens}]_{\text{inf}}^\circ = \underline{B}$ , by (iv) of proposition 2 it follows that  $[\underline{A}, \text{Ens}]_{\text{inf}}$  is inf-complete.

### References

- [1] J. Lambek, *Completions of categories*, Springer Verlag, Berlin, 1966.
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- [3] B. Mitchell, *Theory of categories*, Academic Press, New York, 1965.

Inha Institute of Technology